

KONTROLLTÖÖ

1. Määra semantiliste tabelitega, kas järgmised valemid on üldkehtivad või vääratavad. Vääratavuse korral anna kontramudel.

$$p \wedge (q \vee r) \supset (p \wedge q) \vee r$$

$$\forall x. \exists y. p(x, y) \supset \exists x. \forall y. p(x, y)$$

2. Tõesta loomulikus tuletuses:

$$(p \supset q \vee r) \supset ((q \supset r) \supset (p \vee q \supset r))$$

$$\exists x. p(g(x), a) \supset \exists y. \exists z. p(y, z)$$

3. Tõesta sekventsiarvutuses:

$$\neg(p \wedge \neg r) \wedge p \supset r$$

$$\neg \forall x. \neg p(f(x)) \supset \exists x. p(x)$$

4. Teisenda klauselkujule:

$$\forall y. (\exists x. p(x, y) \supset q(y, z)) \wedge \exists y. (\forall x. r(x, y) \vee q(x, y))$$

5. Leia kõige üldisem unifikaator atomaarvalemitele $p(f(a, c), g(x), x)$, $p(y, g(h(z)), w)$.
6. Resolveeri klauslid $p(a, y) \vee r(y, f(y))$, $\neg r(g(z), w) \vee s(z) \vee \neg t(w, v)$.

LAHENDUSED

1.

$$\mathbf{F}(p \wedge (q \vee r) \supset (p \wedge q) \vee r)$$

$$\mathbf{T}(p \wedge (q \vee r))$$

$$\mathbf{F}((p \wedge q) \vee r)$$

$$\mathbf{T}p$$

$$\mathbf{T}(q \vee r)$$

$$\mathbf{F}(p \wedge q)$$

$$\mathbf{F}r$$

$\mathbf{T}q$		$\mathbf{T}r$
$\mathbf{F}p$	$\mathbf{F}q$	\times
\times	\times	\times

Valem on üldkehtiv.

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$$\mathbf{F}(\forall x. \exists y. p(x, y) \supset \exists x. \forall y. p(x, y))$$

$$\mathbf{T}(\forall x. \exists y. p(x, y))$$

$$\mathbf{F}(\exists x. \forall y. p(x, y))$$

$$\mathbf{T}(\exists y. p(x_0, y))$$

$$\mathbf{F}(\forall y. p(x_0, y))$$

$$\mathbf{T}(p(x_0, x_{01}))$$

$$\mathbf{F}(p(x_0, x_{02}))$$

...

$$\mathbf{T}(\exists y. p(x_s, y))$$

$$\mathbf{F}(\forall y. p(x_s, y))$$

$$\mathbf{T}(p(x_s, x_{s1}))$$

$$\mathbf{F}(p(x_s, x_{s2}))$$

...

○

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Valem on väärata. Tabelist tulenev kontramudel on (D, I) , kus

$D = \{x_0, x_{01}, x_{02}, x_{011}, x_{012}, x_{021}, x_{022}, \dots\}$ ja $I(p)(x_s, x_{s1}) = 1$,

$I(p)(x_s, x_{s2}) = 0$.

On ka lihtsamaid kontramudeleid, juba kahest elemendist põhihulgas piisab.

2.

$$\frac{\frac{\frac{p \supset q \vee r \quad +1 \quad p \quad +4}{q \vee r} \supset \mathcal{E} \quad \frac{q \supset r \quad +2 \quad q \quad +6}{r} \supset \mathcal{E} \quad \frac{+7}{r} \vee \mathcal{E}, -6, -7 \quad \frac{+2 \quad +5}{q \supset r \quad q} \supset \mathcal{E}}{\frac{+3}{p \vee q} \quad \frac{r}{p \vee q \supset r} \supset \mathcal{I}, -3 \quad \frac{+5}{r} \vee \mathcal{E}, -4, -5}}{\frac{p \vee q \supset r \supset \mathcal{I}, -3}{(q \supset r) \supset (p \vee q \supset r)} \supset \mathcal{I}, -2} \supset \mathcal{I}, -1$$

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$$\frac{\frac{\frac{+2}{p(g(x'), a)} \exists \mathcal{I} \quad \frac{+1}{\exists x. p(g(x), a)} \exists \mathcal{I} \quad \frac{\exists z. p(g(x'), z) \quad \exists y. \exists z. p(y, z)}{\exists y. \exists z. p(y, z)} \exists \mathcal{E}, -2, x' \text{ fresh}}{\exists x. p(g(x), a) \supset \exists y. \exists z. p(y, z)} \supset \mathcal{I}, -1}}{\exists x. p(g(x), a) \supset \exists y. \exists z. p(y, z)} \supset \mathcal{I}, -1$$

3.

$$\frac{\frac{\text{Id.} \quad \frac{p, r \rightarrow r}{p \rightarrow \neg r, r} \neg \mathcal{R}}{p \rightarrow p \wedge \neg r, r} \wedge \mathcal{R} \quad \frac{\neg(p \wedge \neg r), p \rightarrow r}{\neg(p \wedge \neg r) \wedge p \rightarrow r} \neg \mathcal{L}}{\rightarrow \neg(p \wedge \neg r) \wedge p \supset r} \supset \mathcal{R}$$

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$$\begin{array}{c}
\text{Id.} \\
\frac{p(f(x')) \rightarrow p(f(x')), \exists x. p(x)}{\frac{p(f(x')) \rightarrow \exists x. p(x)}{\rightarrow \neg p(f(x')), \exists x. p(x)} \exists \mathcal{R}} \exists \mathcal{R} \\
\frac{\frac{p(f(x')) \rightarrow \exists x. p(x)}{\rightarrow \neg p(f(x')), \exists x. p(x)} \neg \mathcal{R}}{\rightarrow \forall x. \neg p(f(x)), \exists x. p(x)} \forall \mathcal{R}, x' \text{ fresh} \\
\frac{\frac{\rightarrow \forall x. \neg p(f(x)), \exists x. p(x)}{\neg \forall x. \neg p(f(x)) \rightarrow \exists x. p(x)} \neg \mathcal{L}}{\rightarrow \neg \forall x. \neg p(f(x)) \supset \exists x. p(x)} \supset \mathcal{R}
\end{array}$$

4.

$$\begin{aligned}
& \forall y. (\exists x. p(x, y) \supset q(y, z)) \wedge \exists y. (\forall x. r(x, y) \vee q(x, y)) \\
& \forall y. (\exists x'. p(x', y) \supset q(y, z)) \wedge \exists y'. (\forall x''. r(x'', y') \vee q(x, y')) \\
& \forall y. (\neg \exists x'. p(x', y) \vee q(y, z)) \wedge \exists y'. (\forall x''. r(x'', y') \vee q(x, y')) \\
& \forall y. (\forall x'. \neg p(x', y) \vee q(y, z)) \wedge \exists y'. (\forall x''. r(x'', y') \vee q(x, y')) \\
& \forall y. \forall x'. (\neg p(x', y) \vee q(y, z)) \wedge \exists y'. \forall x''. (r(x'', y') \vee q(x, y')) \\
& (\forall z. \forall x.) \forall y. \forall x'. \exists y'. \forall x''. ((\neg p(x', y) \vee q(y, z)) \wedge (r(x'', y') \vee q(x, y'))) \\
& ((\neg p(x', y) \vee q(y, z)) \wedge (r(x'', f(z, x, y, x', x'')) \vee q(x, f(z, x, y, x', x''))))
\end{aligned}$$

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5. Kahe atomaarvalemi kõige üldisem unifikaator on $[f(a, c)/y, h(z)/x, h(z)/w]$.
6. Konfliktsete literaalid on A ja $\neg A'$, kus $A = r(y, f(y))$ ja $A' = r(g(z), w)$.
 A ja A' kõige üldisem unifikaator on $[g(z)/y, f(g(z))/w]$.
 Klauslite resolvent on $p(a, g(z)) \vee s(z) \vee \neg t(f(g(z)), v)$.

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