

# Normality, randomness, and the Garden of Eden

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# Introduction

- Cellular automata (CA) are uniform, synchronous model of parallel computation on uniform grids, where the next state of a point is a function of the current state of a finite neighborhood of the point.
- The Garden-of-Eden theorem provides a necessary condition for the global function of a CA in dimension  $d$  to be surjective.
- Also, surjective  $d$ -dimensional CA are **balanced**—every pattern of a given shape has the same number of pre-images.
- Notably, on more complex grids such implications are not respected.
- **Bartholdi's theorem** characterizes **amenable groups** (a class introduced by von Neumann) as those where all surjective CA are balanced.
- We measure the amount by which a surjective CA on a non-amenable group may fail to be balanced.



# The Banach-Tarski paradox (1924)

A closed ball  $U$  in the 3-dimensional Euclidean space can be decomposed into two disjoint subsets  $X, Y$ , both piecewise congruent to  $U$ .

This is due to a series of facts:

- The axiom of choice.
- The group of rotations of the 3-dimensional space has a **free subgroup on two generators**.
- The pieces of the decomposition are not **Lebesgue measurable**.

**What is the role of the group?**



# Amenable groups

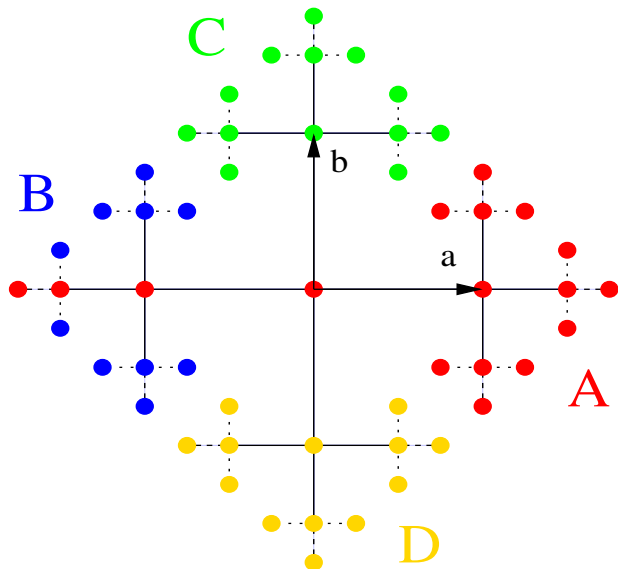
A group  $G$  is **amenable** if there exists a **finitely** additive probability measure  $\mu : \mathcal{P}(G) \rightarrow [0, 1]$  such that:

$$\mu(gA) = \mu(A) \text{ for every } g \in G, A \subseteq G$$

- Subgroups of amenable groups are amenable.
- Quotients of amenable groups are amenable.
- Abelian groups are amenable.
- A group whose **finitely generated** subgroups are **all** amenable, is amenable.



# A paradoxical decomposition of $\mathbb{F}_2$



# Paradoxical groups

A **paradoxical decomposition** of a group  $G$  is a partition  $G = \bigsqcup_{i=1}^n A_i$  such that, for suitable  $\alpha_1, \dots, \alpha_n \in G$ ,

$$G = \bigsqcup_{i=1}^k \alpha_i A_i = \bigsqcup_{i=k+1}^n \alpha_i A_i$$

A **bounded propagation 2:1 compressing map** on  $G$  is a function  $\phi : G \rightarrow G$  such that, for a **finite propagation set**  $S$ ,

- $\phi(g)^{-1}g \in S$  for every  $g \in G$  (bounded propagation) and
- $|\phi^{-1}(g)| = 2$  for every  $g \in G$  (2:1 compression)

A group has a paradoxical decomposition if and only if it has a bounded propagation 2:1 compression map.

Such groups are called **paradoxical**.



# Examples of paradoxical groups

- The free group on two generators is paradoxical.
- Every group with a paradoxical subgroup is paradoxical.
- In particular, every group with a free subgroup on two generators is paradoxical.
- The converse of the previous point is **false!**  
(von Neumann's conjecture; disproved by Ol'shanskii, 1980)
- In fact, there exist paradoxical groups where every element has finite order. (Adian, 1983)



# The Tarski alternative

Let  $G$  be a group. **Exactly one of the following happens:**

- 1  $G$  is amenable.
- 2  $G$  is paradoxical.

**Are there other ways to express that?**





# Cellular automata

A **cellular automaton (CA)** on a group  $G$  is a triple  $\mathcal{A} = \langle Q, \mathcal{N}, f \rangle$  where:

- $Q$  is a finite set of **states**.
- $\mathcal{N} = \{n_1, \dots, n_k\} \subseteq G$  is a finite **neighborhood**.
- $f : Q^k \rightarrow Q$  is a finitary **local function**

The local function induces a **global function**  $F : Q^G \rightarrow Q^G$  via

$$\begin{aligned} F_{\mathcal{A}}(c)(x) &= f(c(x \cdot n_1), \dots, c(x \cdot n_k)) \\ &= f(c^x|_{\mathcal{N}}) \end{aligned}$$

where  $c^x(g) = c(x \cdot g)$  for all  $g \in G$ .

The same rule induces a function over **patterns** with finite **support**:

$$f(p) : E \rightarrow Q, \quad f(p)(x) = f(p^x|_{\mathcal{N}}) \quad \forall p : E\mathcal{N} \rightarrow Q$$



# The Garden-of-Eden theorem

A cellular automaton is **pre-injective** if it satisfies the following condition:

$$\text{if } 0 < |\{g \in G \mid c(g) \neq e(g)\}| < \infty \\ \text{then } F_{\mathcal{A}}(c) \neq F_{\mathcal{A}}(e)$$

**Theorem** (Moore's **Garden-of-Eden theorem**, 1962)

A surjective cellular automaton on  $G = \mathbb{Z}^d$  is pre-injective.

**Theorem** (Myhill, 1963)

A pre-injective cellular automaton on  $G = \mathbb{Z}^d$  is surjective.



# A counterexample on the free group

Let  $G = \mathbb{F}_2$ ,  $Q = \{0, 1\}$ ,  $\mathcal{N} = \{1_G, a, b, a^{-1}, b^{-1}\}$ , and  $f$  the majority rule.

$\mathcal{A}$  is not pre-injective.

- The configuration which has value 1 only on  $1_G$  is updated into the all-0 configuration.

However,  $\mathcal{A}$  is surjective.

- Let  $E \in \mathcal{PF}(G)$  and let  $m = \max\{\|g\| \mid g \in E\}$ .
- Each  $g \in E$  with  $\|g\| = m$  has three neighbors outside  $E$ .
- This allows an argument by induction.



# Prodiscrete topology and product measure

The **prodiscrete topology** of the space  $Q^G$  of configurations is generated by the **cylinders**

$$C(E, p) = \{c : G \rightarrow Q \mid c|_E = p\}$$

The cylinders also generate a  $\sigma$ -algebra  $\Sigma_C$ , on which the **product measure** induced by

$$\mu_{\Pi}(C(E, p)) = |Q|^{-|E|}$$

is well defined.

- $\Sigma_C$  is **not** the Borel  $\sigma$ -algebra unless  $G$  is countable.



# Balancedness

Let  $E$  be a finite nonempty subset of  $G$ ; let  $\mathcal{A} = \langle Q, \mathcal{N}, f \rangle$  be a CA on  $G$ .  $\mathcal{A}$  is  **$E$ -balanced** if for every  $p : E \rightarrow Q$ ,

$$|f^{-1}(p)| = |Q|^{|E\mathcal{N}|-|E|}$$

This is the same as saying that  $\mathcal{A}$  **preserves  $\mu_{\Pi}$** , i.e.,

$$\mu_{\Pi}(F_{\mathcal{A}}^{-1}(U)) = \mu_{\Pi}(U)$$

for every open  $U \in \Sigma_C$ .

**Theorem** (Maruoka and Kimura, 1976)

A CA on  $\mathbb{Z}^d$  is surjective if and only if it is balanced.



# Martin-Löf randomness for infinite words

A **sequential Martin-Löf test** (briefly, M-L test) is a **recursively enumerable**  $U \subseteq \mathbb{N} \times Q^*$  such that the **level sets**  $U_n = \{x \in Q^* \mid (n, x) \in U\}$  satisfy the following conditions:

- 1 For every  $n \geq 1$ ,  $U_{n+1} \subseteq U_n$ .
- 2 For every  $n \geq 1$  and  $m \geq n$ ,  $|U_n \cap Q^m| \leq |Q|^{m-n} / (|Q| - 1)$ .
- 3 For every  $n \geq 1$  and  $x, y \in Q^*$ , if  $x \in U_n$  and  $y \in xQ^*$  then  $y \in U_n$ .

$w \in Q^{\mathbb{N}}$  **fails** a sequential M-L test  $U$  if  $w \in \bigcap_{n \geq 0} U_n Q^{\mathbb{N}}$ .

$w$  is **Martin-Löf random** if  $w$  does not fail any sequential M-L test.

- If  $\eta : \mathbb{N} \rightarrow \mathbb{N}$  is a computable bijection, then  $w$  is M-L random if and only if  $w \circ \eta$  is M-L random.
- It is well known (cf. [Martin-Löf, 1966]) that M-L random words are **normal**.



# What is normality?

Consider the definition for real numbers:

a real number  $x \in [0, 1)$  is **normal in base  $b$**   
if the sequence of its digits in base  $b$  is equidistributed

$x$  is **normal** if it is normal in every base  $b$

A similar definition holds for sequences  $w \in Q^{\mathbb{N}}$ :

- Let  $\text{occ}(u, w) = \{j \geq 0 \mid w_{[j:j+|u|-1]} = u\}$ .
- $w$  is  **$m$ -normal** if for every  $u \in Q^m$ ,

$$\lim_{n \rightarrow \infty} \frac{|\text{occ}(u, w) \cap \{0, \dots, n-1\}|}{n} = |Q|^{-m}$$

**Theorem** (Niven and Zuckerman, 1951)

$w$  is  $m$ -normal over  $Q$  iff it is 1-normal over  $Q^m$ .



# Enumerating the cylinders

Suppose  $G$  is **finitely generated** and has **decidable word problem**.

- Then there is a **computable** bijection  $\phi : \mathbb{N} \rightarrow G$ .
- Also, there is a **computable** function  $m : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$  such that, for all  $i$  and  $j$ , if  $\phi(i) = g$  and  $\phi(j) = h$ , then  $\phi(m(i,j)) = g \cdot h$ .

Then we can enumerate the cylinders as follows:

- First, we enumerate the elementary cylinders:

$$B_{|Q|_{i+j}} = C(g_i, q_j) = \{c : G \rightarrow Q \mid c(\phi(i)) = q_j\}$$

- Next, we define a bijection  $\Psi : \mathcal{PF}(G) \rightarrow \mathbb{N}$  as  $\Psi(X) = \sum_{i \in X} 2^i$  (so that  $\Psi(\emptyset) = 0$ )
- Finally, we enumerate the cylinders as:

$$B'_n = \bigcap_{i \in \Psi^{-1}(n+1)} B_i$$





# Martin-Löf randomness for configurations

Let  $G$  be a f.g. group with decidable word problem.

- We say that  $\mathcal{U}$  is  $\mathcal{V}$ -computable if there exists a r.e.  $A \subseteq \mathbb{N}$  such that

$$U_i = \bigcup_{\pi(i,j) \in A} V_j \quad \forall i \geq 0$$

where  $\pi(i,j) = (i+j)(i+j+1)/2 + j$ .

- A  $B'$ -computable family  $\mathcal{U} = \{U_n\}_{n \geq 0}$  of open subsets of  $Q^G$  is a Martin-Löf  $\mu_{\Pi}$ -test if  $\mu_{\Pi}(U_n) \leq 2^{-n}$  for every  $n \geq 0$ .  
 $c \in Q^G$  fails  $\mathcal{U}$  if  $c \in \bigcap_{n \geq 0} U_n$ .
- $c$  is M-L  $\mu_{\Pi}$ -random if it does not fail any M-L  $\mu_{\Pi}$ -test.



# Two important facts about Martin-Löf randomness

## Theorem (Hertling and Weihrauch)

Let  $\phi : \mathbb{N} \rightarrow G$  an admissible indexing.

$c \in Q^G$  is M-L  $\mu_{\Pi}$ -random if and only if  $c \circ \phi \in Q^{\mathbb{N}}$  is M-L random.

## Theorem (Calude *et al.*, 2001)

Let  $\mathcal{A} = \langle Q, \mathcal{N}, f \rangle$  be a CA on  $\mathbb{Z}^d$ . The following are equivalent:

- 1  $\mathcal{A}$  is surjective.
- 2 For every  $c : \mathbb{Z}^d \rightarrow Q$ , if  $c$  is M-L  $\mu_{\Pi}$  random then so is  $F_{\mathcal{A}}(c)$ .



# Bartholdi's theorem (2010)

Let  $G$  be a group. The following are equivalent:

- 1  $G$  is amenable.
- 2 Every surjective cellular automaton on  $G$  is pre-injective.
- 3 Every surjective cellular automaton on  $G$  preserves the product measure.

**How much does preservation of product measure fail on paradoxical groups?**



# The amount of a failure

**Theorem** (Capobianco, Guillon and Kari)

Let  $G$  be a non-amenable group.

There exist an alphabet  $Q$ , a subset  $U$  of  $Q^G$  such that

$$\mu_{\Pi}(U) = 1,$$

and a surjective cellular automaton  $\mathcal{A}$  over  $G$  with alphabet  $Q$  such that

$$\mu_{\Pi}(F_{\mathcal{A}}^{-1}(U)) = 0.$$



# A surjective, non-balanced CA

Guillon, 2011: improves Bartholdi's counterexample.

Let  $G$  be a non-amenable group,  $\phi$  a bounded propagation 2:1 compressing map with propagation set  $S$ .

Define on  $S$  a total ordering  $\preceq$ .

Define a CA  $\mathcal{A}$  on  $G$  by  $Q = (S \times \{0, 1\} \times S) \sqcup \{q_0\}$ ,  $\mathcal{N} = S$ , and

$$f(u) = \begin{cases} q_0 & \text{if } \exists s \in S \mid u_s = q_0, \\ (p, \alpha, q) & \text{if } \exists (s, t) \in S \times S \mid s \prec t, u_s = (s, \alpha, p), u_t = (t, 1, q), \\ q_0 & \text{otherwise.} \end{cases}$$

Then  $\mathcal{A}$ , although clearly non-balanced, is surjective.

- For  $j \in G$  it is  $j = \phi(js) = \phi(jt)$  for exactly two  $s, t \in S$  with  $s \prec t$ .
- If  $c(j) = q_0$  put  $e(js) = e(jt) = (s, 0, s)$ .
- If  $c(j) = (p, \alpha, q)$  put  $e(js) = (s, \alpha, p)$  and  $e(jt) = (t, 1, q)$ .
- Then  $F(e) = c$ .



# End of the game?

At this point, one might be tempted to reason as such:

- Let  $G$  be a non-amenable group with decidable word problem.
- Let  $c$  be a Martin-Löf random configuration for Guillon's CA.
- There exist some points  $g \in G$  where  $c(g) = q_0$ .
- As  $|S| \geq 2$ ,  $F_{\mathcal{A}}(c)$  cannot have isolated  $q_0$ 's.
- Therefore,  $F_{\mathcal{A}}(c)$  cannot be random.

This argument, albeit convincing, is **wrong**.

- To say that  $F_{\mathcal{A}}(c)$  has no isolated occurrences of  $q_0$ , means that there are some patterns that do not occur in  $F_{\mathcal{A}}(c)$ .
- But  $c$ , being random, is also rich ...
- ... and a rich configuration contains all the preimages of every non-orphan pattern!



## Normality for $d$ -dimensional configurations

It is still sensible to define normality for  $c \in \mathbb{Z}^d$  as follows:

- Let  $E = E(n_1, \dots, n_d) = \prod_{i=1}^d \{0, \dots, n_i - 1\}$ .
- $c : \mathbb{Z}^d \rightarrow Q$  is  $E$ -normal if for every  $p : E \rightarrow Q$ ,

$$\lim_{n \rightarrow \infty} \frac{1}{(2n+1)^d} \cdot |\{x \in \mathbb{Z}^d \mid \|x\| \leq n, c^x|_E = p\}| = \frac{1}{|Q|^{|E|}}$$

But: **why** is this sensible?

- Every  $E$  such as above is a coset for some subgroup of  $\mathbb{Z}^d$ .
- Also, a subgroup of **finite index** of  $\mathbb{Z}^d$  is isomorphic to  $\mathbb{Z}^d$ .

This is **not** true for arbitrary groups!

- If  $G$  is free on two generators, and  $H \leq G$  has index 2, then  $H$  is free on **three** generators!



# So, what is to be done?

The idea:

- Patch the group with patches of a given shape.
- See the state of patches as macrostates.
- Show that  $\mu_{\Pi}$ -almost every configuration is normal with respect to the macrostates.

The problem:

- If we want to fill the group without having the patches overlap, we may be forced to change the underlying group.

The solution: (Kari, 2012)

only patch a portion of the group!





# Normal configurations, modulo some conditions

Let  $G$  be an arbitrary infinite group.

- Let  $E \in \mathcal{PF}(G)$  be nonempty.
- Let  $h : \mathbb{N} \rightarrow G$  be injective.

We define the **lower density**, **upper density**, and **density** of  $U \subseteq G$  according to  $h$ , as the lower limit  $\text{dens inf}_h$ , upper limit  $\text{dens sup}_h$ , and (if exist) limit  $\text{dens}_h$  of

$$\frac{|U \cap h(\{0, \dots, n-1\})|}{n}$$

We say  $c : G \rightarrow Q$  is  **$h$ - $E$ -normal** if for every pattern  $p : E \rightarrow Q$ ,

$$\text{dens}_h \text{occ}(p, c) = |Q|^{-|E|}$$

where  $\text{occ}(p, c) = \{g \in G \mid c^g|_E = p\}$ .



# Sanity check

If  $E \subseteq F$  and  $c$  is  $h$ - $F$ -normal, then it is also  $h$ - $E$ -normal.

- The vice versa is false: for  $h(n) = n, \dots 010101 \dots$  is  $h$ - $\{0\}$ -normal and  $h$ - $\{1\}$ -normal but not  $h$ - $\{0, 1\}$ -normal.

Also, the following are equivalent:

- 1  $c$  is  $h$ - $E$ -normal.
- 2 For every  $p : E \rightarrow Q$ ,  $\text{dens inf}_h \text{occ}(p, c) \geq |Q|^{-|E|}$ .
- 3 For every  $p : E \rightarrow Q$ ,  $\text{dens sup}_h \text{occ}(p, c) \leq |Q|^{-|E|}$ .



## A key lemma

Let  $\mathcal{A} = \langle Q, \mathcal{N}, f \rangle$  be a nontrivial CA on  $G$ .

- Suppose  $\mathcal{A}$  has a spreading state  $q_0$ .
- Let  $s, t$  be two distinct elements of  $\mathcal{N}$ .
- Let  $h : \mathbb{N} \rightarrow G$  be injective.

If  $c : G \rightarrow Q$  is  $h$ - $\{s, t\}$ -normal, then  $F_{\mathcal{A}}(c)$  is **not**  $h$ -1-normal.

- In particular, if  $c$  is  $h$ - $E$ -normal for some  $E \in \mathcal{PF}(G)$  containing  $\mathcal{N}$ , then  $F_{\mathcal{A}}(c)$  is not  $h$ -1-normal.



## The set of non-normal configurations

For  $p : E \rightarrow Q$ ,  $k \geq 1$ , and  $h : \mathbb{N} \rightarrow G$  injective, let

$$L_{h,p,k,n} = \left\{ c : G \rightarrow Q \mid \frac{|\{j < n \mid h(j) \in \text{occ}(p, c)\}|}{n} \leq \frac{1}{|Q|^{|E|}} - \frac{1}{k} \right\}.$$

$\text{dens inf}_h \text{occ}(p, c) < |Q|^{-|E|}$  if and only if there exists  $k \geq 1$  such that

$$c \in \limsup_n L_{h,p,k,n} = \bigcap_{n \geq 1} \bigcup_{m \geq n} L_{h,p,k,m} \stackrel{\text{def}}{=} L_{h,p,k}$$

which is  $\Sigma_C$ -measurable. Then

$$L_{h,E} = \bigcup_{p \in Q^E, k \geq 1} L_{h,p,k}$$

is the set of all the configurations  $c \in Q^G$  that are **not**  $h$ - $E$ -normal.

When is it the case that  $\mu_{\Pi}(L_{h,E}) = 0$ ?



# The Chernoff bound

Let  $Y_0, \dots, Y_{n-1}$  be **independent** nonnegative random variables.

Let  $S_n = Y_0 + \dots + Y_{n-1}$ ,  $\mu = \mu(n) = \mathbb{E}(S_n)$ .

For every  $\delta \in (0, 1)$ ,

$$\mathbb{P}(S_n < \mu \cdot (1 - \delta)) < e^{-\frac{\mu\delta^2}{2}}.$$

In particular, if the  $Y_i$ 's are Bernoulli trials with probability  $p$ , and  $0 < \varepsilon < \min(p, 1 - p)$ , then for  $\delta = \varepsilon/p$

$$\sum_{0 \leq k < n \cdot (p - \varepsilon)} \binom{n}{k} p^k (1 - p)^{n-k} < e^{-\frac{\varepsilon^2 n}{2p}}.$$



# A full set of normal configurations

Suppose that the sets  $h(i)E$ ,  $i \geq 0$ , are pairwise disjoint.

- The random variables

$$Y_i = \left[ c^{h(i)} \Big|_E = p \right]$$

are i.i.d. Bernoulli of parameter  $t = |Q|^{-|E|}$ .

- Set  $S_n = Y_0 + \dots + Y_{n-1}$ . Then for  $\delta = |Q|^{|E|}/k$ ,

$$L_{h,p,k,n} = \{c : G \rightarrow Q \mid S_n < n \cdot |Q|^{-|E|} \cdot (1 - |Q|^{|E|}/k)\}$$

and

$$\mu_{\Pi}(L_{h,p,k,n}) = \mathbb{P}(\{S_n < \mu \cdot (1 - \delta)\}) < e^{-\frac{|Q|^{|E|}}{2k^2} n}$$

- By the [Borel-Cantelli lemma](#), all the  $L_{h,p,k}$  are null sets.

In conclusion:  $\mu_{\Pi}$ -almost every  $c : G \rightarrow Q$  is  $h$ - $E$ -normal



# If it fails, it fails catastrophically

Let  $G$  be a non-amenable group.

- Let  $\mathcal{A} = \langle Q, \mathcal{N}, f \rangle$  be the Guillon CA.
- Let  $E \supseteq \mathcal{N} \cup \{1\}$ .
- Let  $h: \mathbb{N} \rightarrow G$  s.t. the  $h(i)E$ ,  $i \geq 0$ , are pairwise disjoint.
- Then  $\mu_{\Pi}$ -almost every  $c \in Q^G$  is  $h$ - $E$ - and  $h$ -1-normal ...
- ... so none of their preimages can be  $h$ - $E$ -normal!

Hence, the set  $U$  of  $h$ - $E$ -normal configurations satisfies

$$\mu_{\Pi}(U) = 1 \quad \text{and} \quad \mu_{\Pi}(F_{\mathcal{A}}^{-1}(U)) = 0.$$



## Back to randomness

Let  $G$  be an **amenable** group and let  $\mathcal{A} = \langle Q, \mathcal{N}, f \rangle$  be a CA on  $G$ .

- If  $\mathcal{U}$  is  $B'$ -measurable then so is  $F_{\mathcal{A}}^{-1}(\mathcal{U})$ .
- If  $\mathcal{A}$  is surjective and  $\mathcal{U}$  is a M-L  $\mu_{\Pi}$ -test, then so is  $F_{\mathcal{A}}^{-1}(\mathcal{U})$ .
- In these hypotheses, if  $F_{\mathcal{A}}(c)$  fails  $\mathcal{U}$ , then  $c$  fails  $F_{\mathcal{A}}^{-1}(\mathcal{U})$ .

Summarizing:

if  $G$  is amenable,  $\mathcal{A}$  is surjective, and  $c$  is M-L  $\mu_{\Pi}$ -random,  
then  $F_{\mathcal{A}}(c)$  is M-L  $\mu_{\Pi}$ -random





## Fixing a flaw

$a \in Q^{\mathbb{N}}$  is M-L random **relatively to**  $b \in Q^{\mathbb{N}}$  if it is M-L random when computability is considered according to Turing machines with oracle  $b$ .

**Theorem** (van Lambalgen, 1987)

Let  $a, b \in Q^{\mathbb{N}}$  and

$$c(n) = \begin{cases} a(k) & \text{if } n = 2k, \\ b(k) & \text{if } n = 2k + 1. \end{cases}$$

The following are equivalent:

- 1  $c$  is M-L random.
- 2  $a$  is M-L random, and  $b$  is M-L random relative to  $a$ .
- 3  $b$  is M-L random, and  $a$  is M-L random relative to  $b$ .



## Another catastrophic failure!

Let  $G$  be an infinite f.g. group with decidable word problem.  
For every nonempty  $E \in \mathcal{PF}(G)$  there exists a computable injective function  $h : \mathbb{N} \rightarrow G$  such that:

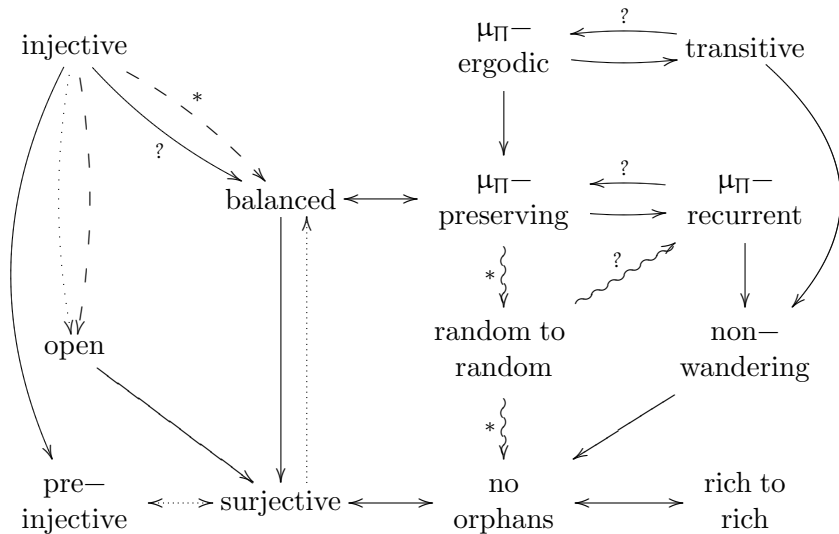
- 1  $h(\mathbb{N})$  is a recursive subset of  $G$  with infinite complement.
- 2  $h(n)E \cap h(m)E = \emptyset$  for every  $n \neq m$ .
- 3 For any alphabet  $Q$ , every M-L  $\mu_{\Pi}$ -random configuration  $c : G \rightarrow Q$  is  $h$ - $E$ -normal. (This follows from van Lambalgen's theorem.)

Let then  $\mathcal{A}$  be the Guillon CA.

- Construct  $h$  as above with  $E = \mathcal{N} \cup \{1\}$ .
- Let  $c : G \rightarrow Q$  be a M-L  $\mu_{\Pi}$ -random configuration.
- Because of the above lemma,  $F_{\mathcal{A}}(c)$  cannot be random.
- For the same reason, none of the preimages of  $c$  can be random.



# A diagram of implications



## Conclusions and future work

- The characterizations of surjective CA listed in [Calude et al., 2001] actually hold on arbitrary amenable groups—and precisely on those.
- Among those, preservation of the product measure is the one that fails **catastrophically** on paradoxical groups.
- Does **Myhill's theorem** fail for paradoxical groups?  
(This problem seems **very** difficult!)
- Are there injective CA which are not balanced?  
(If no such CA exists, then **Gottschalk's conjecture** is true.)
- Does there exist a CA that sends a nonrich configuration into a rich one?

# Thank you for attention!

Any questions?

