Diagrammatic sets between rewriting and topology

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arXiv:1909.07639

Representable diagrammatic sets as a model of weak higher categories.

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Dimension 0: (labelled) abstract rewrite system

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Dimension 1: string rewriting system

Dimension 0: (labelled) abstract rewrite system



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Dimension 1: string rewriting system



Dimension 2: algebraic theories / monoidal categories



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Rewriting objects of dimension > 2:

- higher/homotopical algebra
- *k*-manifolds in *n*-space (e.g. k=1, n=3: knots and braids)

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Higher-dimensional "rewrites" for a fixed base dimension:confluence, coherence...

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Except cells have a direction ("Computationally aware" homotopy theory?)

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"Point-set" cells not available — need a combinatorial notion of 1. **directed cells** 2. their **pasting**

Directed *n*-cells are modelled by *n*-globes, the objects classifying *n*-cells in a strict ω-category

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Problem:

The pasting maps are **not sound** for the interpretation of "rewrite systems as CW complexes"

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Problem:

The pasting maps are **not sound** for the interpretation of "rewrite systems as CW complexes"

...plus other technical issues

Towards diagrammatic sets

Let the CW complex interpretation guide the choice of a framework

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We can associate to a CW complex its face poset...



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We can associate to a CW complex its face poset...



and to a pasting diagram its oriented face poset.

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Regular CW complex X: pasting maps are homeomorphisms with their image

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A classical theorem of combinatorial topology

A regular CW complex is specified up to cellular homeomorphism by its face poset



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Conjecture

A regular pasting diagram is specified up to cellular isomorphism by its oriented face poset

 Directed *n*-cells are modelled by regular directed complexes (which are oriented face posets of regular pasting diagrams)
with a greatest element of rank *n* (so the underlying poset is the face poset of a regular CW *n*-ball)

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 Pasting is given by maps of posets that are compatible functorially with both realisations

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Maps factor into

■ injections, giving face operations

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 - \rightarrow sub-diagrams, substitutions in context

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 \rightarrow "nullary" operations in universal algebra

Enough for higher-dimensional rewriting?

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Good geometric realisation; can be used to construct CW complexes

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(and I think these are important in higher algebra)

Good geometric realisation; can be used to construct CW complexes

(in fact, diagrammatic sets satisfy a version of the *homotopy hypothesis* — one can reason about spaces/homotopy types in terms of their diagrammatic nerve, as with simplicial sets)



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Diagrammatic sets have *some* features of a model (pasting, units)

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Idea: higher categories \rightarrow diagrammatic sets with an **internal** notion of *weak composition*

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Idea: higher categories \rightarrow diagrammatic sets with an **internal** notion of *weak composition*

(in the spirit of categorical semantics: syntax and semantics in the same universe)

Computational meaning of composition:

A diagram x can be **substituted** in every context with a cell [x]

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A diagram x can be substituted in every context with another diagram y (and vice versa)

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There are special equivalence cells $x \Rightarrow y$, $y \Rightarrow x$, which mediate between all cells containing x and all cells containing y in their boundary





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And this equivalence should be witnessed by **3-dimensional** equivalence cells...



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whose definition involves 4-dimensional equivalence cells, etc

This is a **coinductive** definition.

Let X be a diagrammatic set. For all subsets $A \subseteq \operatorname{Cell}(X)$, define

$$\begin{split} \mathcal{F}(A) &:= \{ x : U \to X \mid \text{ for all } \alpha \in \{+, -\} \text{ and} \\ (\Lambda \hookrightarrow W, \lambda : \Lambda \to X) \in \mathcal{D}iv(x, \partial^{\alpha}U), \\ \text{ there exists } (h : W \to X) \in A \text{ such that } h|_{\Lambda} = \lambda \}; \end{split}$$

Then \mathcal{F} is an order-preserving map on $\mathcal{P}(\operatorname{Cell}(X))$. Its **greatest** fixed point is the set $\mathcal{E}qX$ of equivalence cells of X.

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Then \mathcal{F} is an order-preserving map on $\mathcal{P}(\operatorname{Cell}(X))$. Its **greatest** fixed point is the set $\mathcal{E}qX$ of equivalence cells of X.

Proof method: if $A \subseteq \mathcal{F}(A)$, then $A \subseteq \mathcal{E}qX$.

Representable diagrammatic set (RDS)

A diagrammatic set where, for all diagrams x, there exist cells [x], [x]' and equivalence cells $x \Rightarrow [x], [x]' \Rightarrow x$.

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A diagrammatic set where, for all diagrams x, there exist cells [x], [x]' and equivalence cells $x \Rightarrow [x], [x]' \Rightarrow x$.

For all diagrams x and y, let $x \simeq y$ if there exists an equivalence $x \Rightarrow y$.

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Representable diagrammatic set (RDS)

A diagrammatic set where, for all diagrams x, there exist cells [x], [x]' and equivalence cells $x \Rightarrow [x], [x]' \Rightarrow x$.

For all diagrams x and y, let $x \simeq y$ if there exists an equivalence $x \Rightarrow y$.

Theorem

In a representable diagrammatic set,

- 1 all degenerate cells are equivalence cells, and
- 2 \simeq is an equivalence relation.

- "Groupoidal" RDSs (in which every cell is an equivalence) model all homotopy types.
- 2 Conditional to the conjecture on regular pasting diagrams, strict ω-categories embed as a full subcategory (if one takes morphisms that preserve a choice of weak composites)
- 3 There are *n*-truncated RDSs corresponding to weak *n*-categories. 2-truncated RDSs are equivalent to bicategories