# ITT8040 — Cellular Automata Lecture 8

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- Nilpotency.
- Surjectivity in dimension 2 (or greater).
- Reversibility in dimension 2 (or greater).

How can we ensure reversibility in CA?

A lattice gas cellular automaton (LGCA) is a CA (S, d, N, f) where:

- ► The set of states S = S<sub>1</sub> × ... × S<sub>m</sub> has as many channels as the elements of N = {n<sub>1</sub>,..., n<sub>m</sub>}.
- ▶ The global evolution function is a sequence of two steps:
  - 1. Propagation: each track  $S_i$  is shifted by  $\vec{n}_i$ ;
  - 2. Interaction:

a transformation  $\pi: S \to S$  is performed.

Hardy, de Pazzis and Pomeau

- Square grid, four directions.
- Particle moving along channels between nodes of the grid.
- Propagation: to nearest neighbor.
- Interaction:
  - If exactly two particles arrive from opposite directions, then they bounce at a right angle.
  - In all other cases, the particles go straight ahead.









HPP has several problems, which make it unsuited as a gas model:

- Lattice directions are privileged.
- Horizontal momentum is preserved along horizontal lines.
   Such conservation law does not hold for real gasses.
- Frisch, Hasslacher and and Pomeau
  - Triangular grid, six directions.
  - Particle moving along channels between nodes of the grid.
  - Propagation: to nearest neighbor.
  - Interaction:
    - If exactly two particles arrive from opposite directions, then they bounce by 60 degrees at a random direction.
    - If three particles arrive by 120 degrees from each other, then they bounce away.

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In all other cases, the particles go straight ahead.





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The following are equivalent:

- 1. The LGCA is reversible.
- 2. The global function of the LGCA interaction phase is bijective.
- 3. The local function of the LGCA interaction phase is a permutation.

As a side effect:

if a LGCA is not injective, then it is not surjective either

- A partitioned cellular automaton is a CA (S, d, N, f) where:
  - ► The set of states S = S<sub>1</sub> × ... × S<sub>m</sub> has as many tracks as the elements of N = {n
    <sub>1</sub>,...,n
    <sub>m</sub>}.
  - ► The global evolution function is a sequence of two steps:
    - 1. first, each track  $S_i$  is shifted by  $\vec{n}_i$ ;
    - 2. next, a permutation  $\pi: S \to S$  is performed;

that is, G is a composition of partial shifts and a point-based permutation.

#### Theorem

There exists an algorithm that, given in input an arbitrary Turing machine, returns a reversible one-dimensional partitioned cellular automaton that simulates the Turing machine in real time.

#### Corollary

There exists a reversible 1D PCA and a subset F of its states for which the following problem is r.e.-complete:

given a finite configuration c, determine whether it ultimately evolves into a configuration which has some of its states in F. Reversibility means that no information is erased.

• We need a "garbage track" to take care of previous states. Let the Turing machine  $M = (Q, q_I, q_A, q_R, \Gamma, \Sigma, \delta)$  be given. Set  $S = S_1 \times S_2 \times S_3 \times S_4$  with:

- $S_1 = \Gamma$ .
- $\blacktriangleright S_2 = Q \sqcup \{0\}.$
- $\triangleright \ S_3 = Q \sqcup \{0\}.$

Until now, the construction is the same as for the non-reversible case, except that we use two tracks to encode motion of the read-write head.

 $\triangleright \ S_4 = (Q \times \Gamma \times \{-1, +1\}) \sqcup \{0\}.$ 

This is the "garbage track", which will be moved two slots per time unit, so that it does not interfere with the computation.

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## Simulating TM by RPCA: the local function

Set  $N = \{0, -1, +1, +2\}$ .

- Track 1 does not move.
- Track 2 is shifted one position to the left.
- Track 3 is shifted one position to the right.
- Track 4 is shifted two positions to the right.

Define  $\pi: S \to S$  as a permutation that fulfills the conditions:

• if 
$$\delta(q, a) = (q', a', -1)$$
  
then  $(a, q, 0, 0) \mapsto (a', q', 0, (q, a, -1))$   
and  $(a, 0, q, 0) \mapsto (a', q', 0, (q, a, +1));$ 

• if 
$$\delta(q, a) = (q', a', +1)$$
  
then  $(a, q, 0, 0) \mapsto (a', 0, q', (q, a, -1))$   
and  $(a, 0, q, 0) \mapsto (a', 0, q', (q, a, +1));$ 

•  $(a, 0, 0, g) \mapsto (a, 0, 0, g)$  whatever g is.

### Kari, 1996

- No additional state
- Proved in dimension 1 and 2
- Conjectured for higher dimension

Durand-Lose, 1999

- Additional state
- Worsk in arbitrary dimension

Let (S, 1, N, f) be a one-dimensional CA. The *m*-block presentation of the CA is determined by the block merging function

$$B_m: S^{\mathbb{Z}} \to (S^m) \to \mathbb{Z}$$
 such that  $B_m(c)(i) = c_{[mi,mi+m-1]}$ 

and the block splitting function

 $B_m^{-1}:(S^m)\to\mathbb{Z}\to S^{\mathbb{Z}}$  such that  $B_m^{-1}(e)(i)=e(\lfloor i/m 
floor)(i \mod m)$ 

so that the global function of the m-th higher block presentation is

$$B_m \circ G \circ B_m^{-1}$$

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A similar idea works in higher dimension.

### Reversible 1D CA as PCA: construction

Let G be the global function of a reversible 1D CA. Suppose G and  $G^{-1}$  are both defined by a radius-r neighborhood. Set n = 3r. Define the set of right stairs:

$$\mathcal{R} = \{(c_{[0,2r-1]}, G(c)_{[-r,r-1]}) \mid c \in S^{\mathbb{Z}}\} \subseteq^{2r} \times S^{2r}$$

and the set of left stairs:

$$\mathcal{L} = \{(\mathcal{G}(c)_{[0,2r-1]}, c_{[-r,r-1]}) \mid c \in S^{\mathbb{Z}}\} \subseteq S^{2r} \times S^{2r}$$

The function  $\varphi: \mathcal{S}^{6r} \to \mathcal{R} \times \mathcal{L}$  defined by

$$\phi(c_0,\ldots,c_{6r-1}) = ((c_{[4r,6r-1]},G(c)_{[3r,5r-1]}),(G(c)_{[r,3r-1]},c_{[0,2r-1]}))$$

is a bijection, and so is  $\psi:\mathcal{R}\times\mathcal{L}\to S^{6r}$  defined by

$$\psi((c_{[4r,6r-1]}, G(c)_{[3r,5r-1]}), (G(c)_{[r,3r-1]}, c_{[0,2r-1]})) = G(c)_{[0,6r-1]}$$

Consider the radius-1/2 PCA with set of states  $\mathcal{R}\times\mathcal{L}$  and permutation function

 $\pi=\psi\circ\varphi$ 

Such a PCA is isomorphic to the 6*r*-block presentation of  $G \circ \sigma^{3r}$ .