# ITT8040 Cellular Automata Solutions to Assignment 5 

## Exercise 1

Let $A=(S, 1, N, f)$ be a one-dimensional CA vith global function $G$. Consider the de Bruijn labeled graph $\Gamma=(V, E)$ of $A$. If a configuration $c$ is represented by the labeling of a bi-infinite path $\pi$, then the sequence of the nodes touched by $\pi$ represents, up to a fixed shift, a configuration $e$ such that $G(e)=c$.

## (a) Every non-GoE spatially periodic point has a periodic preimage

Let $p$ be the period of $c$, that is, let $p>0$ and $c(n+p)=c(n)$ for every $n \in \mathbb{Z}$. Consider a bi-infinite path $\pi$ in $\Gamma$ whose labeling represents $c$ : as $\Gamma$ is finite, there must be a sequence $s$ of $p$ consecutive edges in $\Gamma$ that occurs infinitely often in $\pi$. The nodes touched between two consecutive occurrences of $s$ define a periodic configuration which, up to a translation, is a preimage of $c$.

## (b) If $A$ has a fixed point then it also has a spatially periodic fixed point

Let $m$ be the neighborhood range of the CA. Let $\pi$ be a bi-infinite path in $\Gamma$ whose labeling represent a fixed point $c$. Divide $\pi$ in slices of length $m$ : as $\pi$ is infinite and the finite paths of $m$ edges in $\Gamma$ are finitely many, there must exist a slice $s$ which is repeated infinitely often in $\pi$. The slice of $\pi$ between two consecutive occurrences of $s$ is easily seen to determine a periodic fixed point.

## Exercise 2

Modify the semi-algorithm $S$ into an algorithm $S^{\prime}$ that operates as follows:

1. First, use $A$ to compute the value $f(x)$.
2. Store the value $f(x)+1$ at a location $n$.
3. Then, reproduce the behavior of $S$ on $x$, but decrease the value of $n$ by one unit each time a step from $S$ is performed.
4. If the answer "yes" is obtained, return "yes".
5. If the answer "no" is obtained, return "no".

6 . If the value of $n$ reaches zero, return "no".
If $x$ is a "yes" instance of $P$, then $S$ returns "yes" in at most $f(x)$ steps, so $S^{\prime}$ halts and returns "yes". If $x$ is a "no" instance of $P$, then either $S$ ultimately halts and returns "no", or it runs for no less than $f(x)+1$ steps, so $S^{\prime}$ halts and returns "no".

