

# Concrete Mathematics

## Exercises from Chapter 7

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### Warmups

#### Exercise 7.7

Solve the recurrence:

$$\begin{aligned}g_0 &= 1 \\g_n &= g_{n-1} + 2g_{n-2} + \dots + ng_0\end{aligned}$$

*Solution.* Let  $G(z)$  be the generating function of the sequence  $\langle g_0, g_1, g_2, \dots \rangle$ . The recurrence above tells us that  $G(z)$  is the convolution of itself with the generating function of the sequence  $\langle 0, 1, 2, \dots \rangle$ , which is  $z/(1-z)^2$ : except for the first term, which is 1 instead of  $0 = 0 \cdot g_0$ . Hence,

$$G(z) = 1 + \frac{zG(z)}{(1-z)^2}$$

which rewrites as

$$(1-z)^2G(z) = (1-z)^2 + zG(z)$$

which yields

$$(1-3z+z^2)G(z) = (1-z)^2$$

that is,

$$G(z) = \frac{1-2z+z^2}{1-3z+z^2} = 1 + \frac{z}{1-3z+z^2}$$

The first summand on the right-hand side is clearly the generating function of the sequence  $a_n = [n=0]$ ; the second one, is the generating function of  $b_n = f_{2n}$ , where  $f_n$  is the  $n$ th Fibonacci number. Therefore,  $g_n = f_{2n} + [n=0]$ .

## An exercise by Albert R. Meyer and Ronitt Rubinfeld

Let  $a_n$  be the number of string on a ternary alphabet that contain a double character, *i.e.*, a sequence  $xx$  with  $x$  a letter.

1. Find a recurrence for  $a_n$ .
2. Let  $G(z)$  be the generating function of the sequence  $\langle a_0, a_1, a_2, \dots \rangle$ . Prove that

$$G(z) = \frac{-z}{1-2z} + \frac{z}{(1-2z)(1-3z)} \quad (1)$$

3. Find  $r$  and  $s$  such that

$$\frac{1}{(1-2z)(1-3z)} = \frac{r}{1-2z} + \frac{s}{1-3z} \quad (2)$$

4. Find a closed form for  $a_n$ .

*Solution.*

**Point 1.** Call “good” a sequence with a double letter, “bad” a sequence without. Then  $a_n$  is the number of good sequences of  $n$  letters.

There are no good sequences of length 0 or 1, so  $a_0 = a_1 = 0$ . For  $n \geq 2$ , a good sequence of length  $n$  can be obtained from a sequence of length  $n-1$  in two ways: Either we add any letter at the end of a good sequence, or we duplicate the last letter of a bad sequence. No string allows applying both methods, so the number of good strings of length  $n$  is three times the number of good strings of length  $n-1$ , plus the number of bad strings of length  $n-1$ . Thus,

$$a_n = 3a_{n-1} + (3^{n-1} - a_{n-1}) = 2a_{n-1} + 3^{n-1} \quad (3)$$

for every  $n \geq 2$ .

**Point 2.** We want to rewrite (3) in terms of generating functions. Since the recurrence only holds for  $n \geq 2$ , we must consider formal power series whose constant and linear term are zero. Recall that  $1/(1-\alpha z)$  is the generating function of the sequence of the powers of  $\alpha$ .

Let  $G(z)$  be the generating function of  $\langle a_0, a_1, a_2, a_3, \dots \rangle$ . Then  $zG(z)$  is the generating function of  $\langle 0, a_0, a_1, a_2, \dots \rangle$ , while that of  $\langle 0, 1, 3, 9, 27, \dots \rangle$  is  $z/(1-3z)$ . Therefore, in terms of generating functions, (3) is rewritten as

$$G(z) - a_0 - a_1z = 2z(G(z) - a_0) + z \left( \frac{1}{1-3z} - 1 \right) :$$

which, since  $a_0 = a_1 = 0$ , gives

$$(1 - 2z)G(z) = \frac{z}{1 - 3z} - z$$

which yields (1).

**Point 3.** Equation (2) is satisfied if and only if  $r \cdot (1 - 3z) + s \cdot (1 - 2z) = 1$  whatever  $z$  is. For  $z = 1/2$  we find  $r \cdot (1 - 3/2) = 1$ , so  $r = -2$ . For  $z = 1/3$  we find  $s \cdot (1 - 2/3) = 1$ , so  $s = 3$ . In conclusion,

$$\frac{1}{(1 - 2z)(1 - 3z)} = \frac{3}{1 - 3z} - \frac{2}{1 - 2z}.$$

**Point 4.** We can rewrite (1) as follows:

$$\begin{aligned} G(z) &= \frac{-z(1 - 3z) + z}{(1 - 2z)(1 - 3z)} \\ &= \frac{-z + 3z^2 + z}{(1 - 2z)(1 - 3z)} \\ &= 3z^2 \cdot \left( \frac{3}{1 - 3z} - \frac{2}{1 - 2z} \right) \end{aligned}$$

For  $n \geq 2$  it must then be:

$$\begin{aligned} a_n &= [z^n]G(z) \\ &= 3 \cdot \left( [z^{n-2}] \left( \frac{3}{1 - 3z} \right) - [z^{n-2}] \left( \frac{2}{1 - 2z} \right) \right) \\ &= 9 \cdot 3^{n-2} - 6 \cdot 2^{n-2} \\ &= 3 \cdot (3^{n-1} - 2^{n-1}). \end{aligned}$$

### Exercise 7.11

Let  $a_n = b_n = c_n = 0$  for  $n < 0$ , and

$$A(z) = \sum_n a_n z^n ; \quad B(z) = \sum_n b_n z^n ; \quad C(z) = \sum_n c_n z^n$$

1. Express  $C$  in terms of  $A$  and  $B$  when  $c_n = \sum_{j+2k \leq n} a_j b_k$ .
2. Express  $A$  in terms of  $B$  when  $nb_n = \sum_{k=0}^n 2^k a_k / (n - k)!$

3. Express  $A$  in terms of  $B$  when  $a_n = \sum_{k=0}^n \binom{r+k}{k} b_{n-k}$ . Then, construct  $\{f_n(r)\}_{n \geq 0}$  such that  $b_n = \sum_{k=0}^n f_k(r) a_{n-k}$ .

*Solution.*

**Point 1.** We know that, if  $a_n = [z^n]G(z)$ , then  $\sum_{k \leq n} a_k = [z^n] \frac{G(z)}{1-z}$ . Then we can solve point 1 as soon as we find  $G(z)$  such that  $[z^n]G(z) = \sum_{j+2k=n} a_j b_k$ . But the latter is the coefficient of index  $n$  of the convolution of  $A$  with a power series whose odd-indexed coefficients are 0, and whose coefficient of index  $2k$  is  $b_k$ : such function is precisely  $B(z^2)$ . Therefore,

$$C(z) = \frac{A(z)B(z^2)}{1-z}$$

We will leave the next points for the next lecture.