Asymptotics ITT9132 Concrete Mathematics Lecture 16 – 13 May 2019

Chapter Nine O Manipulation Two Asymptotic Tricks Euler's Summation Formula Final Summations



### Contents

### 1 O Manipulation

- Z Two Asymptotic Tricks
   Trick 1: Bootstrapping
   Trick 2: Trading tails
- 3 Euler's Summation Formula
- 4 Final Summations
  - A bell-shaped summand
  - Stirling's approximation



### Next section

### **1** O Manipulation

- Two Asymptotic Tricks
  Trick 1: Bootstrapping
  Trick 2: Trading tails
- 3 Euler's Summation Formula
- 4 Final Summations
  A bell-shaped summand
  Stirling's approximation



#### Theorem

Let 
$$G(z) = \sum_{k \ge 0} g_k z^k$$
 have convergence radius  $R > 0$ . If either

•  $f(n) \prec 1$ , or

• 
$$R = \infty$$
 and  $f(n) = O(1)$ ,

then for every  $m \ge 1$ , if h(n) = O(f(n)) then:

$$G(h(n)) = \sum_{0 \leq k < m} g_k(h(n))^k + O((f(n))^m)$$

#### Remark:

- $f(n) \prec 1$  if and only if  $\lim_{n\to\infty} f(n) = 0$ : that is, f vanishes at infinity.
- f(n) = O(1) if and only if  $|f(n)| \leq C$  for some C > 0 and every n: that is, f is bounded.
- The thesis is equivalent to saying that:

$$\sum_{k \ge m} g_k(h(n))^k = O((f(n))^m) \text{ whenever } h(n) = O(f(n)).$$



#### Theorem

Let  $G(z) = \sum_{k \ge 0} g_k z^k$  have convergence radius R > 0. If either:

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$$R = \infty$$
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$$G(h(n)) = \sum_{0 \leq k < m} g_k(h(n))^k + O((f(n))^m)$$

Examples:

In 
$$(1 + O(1/n)) = O(1/n)$$
.
In  $(1 + 1/n) = \frac{1}{n} - \frac{1}{2n^2} + \frac{1}{3n^3} + O\left(\frac{1}{n^4}\right)$ .
 $e^{\frac{\ln n}{n}} = 1 + \frac{\ln n}{n} + \frac{(\ln n)^2}{2n^2} + O\left(\left(\frac{\ln n}{n}\right)^3\right)$ .
 $e^{1-1/n} = \frac{5}{2} - \frac{2}{n} + \frac{1}{2n^2} + O\left(\left(1 - \frac{1}{n}\right)^3\right)$ .
This holds because  $e^z = \sum_{k \ge 0} \frac{z^k}{k!}$  has infinite convergence radius.



#### Theorem

Let  $G(z) = \sum_{k \ge 0} g_k z^k$  have convergence radius R > 0. If either:

- $f(n) \prec 1$ , or
- $R = \infty$  and f(n) = O(1),

then for every  $m \ge 1$ , if h(n) = O(f(n)) then:

$$G(h(n)) = \sum_{0 \leq k < m} g_k(h(n))^k + O((f(n))^m)$$

Proof for  $R < \infty$  and  $f \prec 1$ :

- Let h(n) = O(f(n)). Choose C > 0 such that  $|h(n)| \leq C \cdot |f(n)|$  for every  $n \ge 0$ .
- Fix  $\delta \in (0,R)$  Then  $K_m = C^m \cdot \sum_{k \geqslant m} |g_k| \delta^{k-m} < \infty$
- Choose  $n_0$  such that  $|f(n)| < \delta/C$  for every  $n \ge n_0$ : for such n,  $|h(n)| < \delta$ .
- Then for  $n \ge n_0$ :

$$\left|\sum_{k \geqslant m} g_k(h(n))^k\right| < |h(n)|^m \cdot \sum_{k \geqslant m} |g_k| \cdot \delta^{k-m} \leqslant K_m \cdot |f(n)|^m$$



#### Theorem

Let  $G(z) = \sum_{k \ge 0} g_k z^k$  have convergence radius R > 0. If either.

- $f(n) \prec 1$ , or
- $R = \infty$  and f(n) = O(1),

then for every  $m \ge 1$ , if h(n) = O(f(n)) then:

$$G(h(n)) = \sum_{0 \leq k < m} g_k(h(n))^k + O((f(n))^m)$$

Proof for  $R = \infty$  and f(n) = O(1):

- Let h(n) = O(f(n)). Choose C > 0 such that  $|h(n)| \leq C \cdot |f(n)|$  for every  $n \ge 0$ .
- Fix  $\delta > 0$  such that  $|f(n)| \leq \delta/C$  for every *n*. Then  $|h(n)| \leq \delta$  for every *n* too.
- As  $R = \infty$ ,  $K_m = C^m \cdot \sum_{k \ge m} |g_k| \delta^{k-m} < \infty$ .
- Then for every n:

$$\left|\sum_{k \geq m} g_k(h(n))^k\right| \leq |h(n)|^m \cdot \sum_{k \geq m} |g_k| \cdot \delta^{k-m} \leq K_m \cdot |f(n)|^m$$



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### 1 O Manipulation

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 Trick 1: Bootstrapping
 Trick 2: Trading tails

3 Euler's Summation Formula

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### 1 O Manipulation

- Z Two Asymptotic Tricks
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   Trick 2: Trading tails
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### Bootstrapping

To get an asymptotics for the solution of a recurrence:

- 1 first, take a rough estimate;
- 2 next, substitute the estimate into the recurrence.

This is an example of "pulling oneself up by one's bootstraps".

#### Problem

Find an asymptotic estimate for  $g_n = [z^n]G(z)$  where  $G(z) = e^{\sum_{k \geqslant 1} rac{z^k}{k^2}}$  .



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#### Step 1: Look for a simple asymptotics

As we have an exponential, it might be a good idea to differentiate by series, and see if a more manageable formula appears:

$$\sum_{n \ge 1} ng_n z^{n-1} = G(z) \cdot \sum_{k \ge 1} \frac{z^{k-1}}{k} = \sum_{n \ge 1} \sum_{0 \le k < n} \frac{g_k}{n-k} z^{n-1}$$

which gives us the recurrence:

$$ng_n = \sum_{k=0}^{n-1} \frac{g_k}{n-k} \quad \forall n \ge 1.$$



#### Problem

Find an asymptotic estimate for  $g_n = [z^n]G(z)$  where  $G(z) = e^{\sum_{k \ge 1} \frac{z^k}{k^2}}$ .

#### Step 1: Look for a simple asymptotics

A numerical computation of the first values gives:

п	0	1	2	3	4	5	6
gn	1	1	$\frac{3}{4}$	$\frac{19}{36}$	$\frac{107}{288}$	$\frac{641}{2400}$	$\frac{51103}{259200}$

It looks like all the terms are positive and not larger than 1: which can easily be proved by induction. Then:  $g_n = O(1)$ .



#### Problem

Find an asymptotic estimate for  $g_n=[z^n]G(z)$  where  $G(z)=e^{\sum_{k\geqslant 1}rac{z^k}{k^2}}$  .

#### Step 2: Plug in the asymptotics

If we replace the  $g_k$  on the right-hand side with O(1), we find:

$$ng_n = \sum_{k=0}^{n-1} \frac{O(1)}{n-k} = \sum_{k=1}^n \frac{O(1)}{k} = H_n \cdot O(1) = O(\log n).$$

Then:  $g_n = O\left(\frac{\log n}{n}\right)$ .



#### Problem

Find an asymptotic estimate for 
$$g_n = [z^n]G(z)$$
 where  $G(z) = e^{\sum_{k \ge 1} \frac{z^n}{k^2}}$ 

#### Step 2: Plug in the asymptotics again

If we replace the  $g_k$  on the right-hand side with  $O(\log(k)/k)$  and put the first summand out of the sum, we find:

$$ng_n = \frac{1}{n} + \sum_{k=1}^{n-1} \frac{1}{n-k} O\left(\frac{\log k}{k}\right) = \frac{1}{n} + \sum_{k=1}^{n-1} \frac{O(\log k)}{k(n-k)} = \frac{1}{n} + \frac{2}{n} H_{n-1} O(\log n),$$

because 
$$\frac{1}{k(n-k)} = \frac{1}{n} \left( \frac{1}{k} + \frac{1}{n-k} \right)$$
 and  $k = O(n)$ . Then  $g_n = O\left( \left( \frac{\log n}{n} \right)^2 \right)$ 



#### Problem

Find an asymptotic estimate for  $g_n = [z^n]G(z)$  where  $G(z) = e^{\sum_{k \ge 1} \frac{z^k}{k^2}}$ .

#### Step 3: Manipulate the estimate obtained

We now rewrite our original recurrence by "pulling out the largest part":

$$ng_n = \sum_{k=0}^{n-1} \frac{g_k}{n} + \sum_{k=0}^{n-1} g_k \left(\frac{1}{n-k} - \frac{1}{n}\right) = \frac{1}{n} \sum_{k \ge 0} g_k - \frac{1}{n} \sum_{k \ge n} g_k + \frac{1}{n} \sum_{k=0}^{n-1} \frac{kg_k}{n-k}$$



#### Problem

Find an asymptotic estimate for  $g_n = [z^n]G(z)$  where  $G(z) = e^{\sum_{k \geqslant 1} \frac{z^k}{k^2}}$ .

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As  $\sum_{k \ge 1} \frac{z^k}{k^2}$  converges at z = 1, the first sum is "simply"  $G(1) = e^{\sum_{k \ge 1} 1/k^2} = e^{\pi^2/6}$ . For the second sum, we observe that:

$$\sum_{k>n} \frac{\log^2 k}{k^2} < \sum_{m \ge 1} \sum_{n^m < k \le n^{m+1}} \frac{\log^2 n^{m+1}}{k(k-1)} < \sum_{m \ge 1} \frac{(m+1)^2 \log^2 n}{n^m} = O\left(\frac{\log^2 n}{n}\right),$$

so 
$$\sum_{k \ge n} g_k = \sum_{k \ge n} O\left(\frac{\log^2 k}{k^2}\right) = O\left(\sum_{k \ge n} \frac{\log^2 k}{k^2}\right) = O\left(O\left(\frac{\log^2 n}{n}\right)\right) = O\left(\frac{\log^2 n}{n}\right).$$

#### Problem

Find an asymptotic estimate for  $g_n = [z^n]G(z)$  where  $G(z) = e^{\sum_{k \ge 1} \frac{z^k}{k^2}}$ .

#### Step 3: Manipulate the estimate obtained

We now rewrite our original recurrence by "pulling out the largest part":

$$ng_n = \sum_{k=0}^{n-1} \frac{g_k}{n} + \sum_{k=0}^{n-1} g_k \left(\frac{1}{n-k} - \frac{1}{n}\right) = \frac{1}{n} \sum_{k \ge 0} g_k - \frac{1}{n} \sum_{k \ge n} g_k + \frac{1}{n} \sum_{k=0}^{n-1} \frac{kg_k}{n-k}$$

For the third and last sum, we again plug in the asymptotics:

$$\sum_{0 \le k < n} \frac{kg_k}{n-k} = \sum_{0 < k < n} \frac{k}{n-k} O\left(\left(\frac{\log k}{k}\right)^2\right) = O\left(\sum_{0 < k < n} \frac{\log^2 k}{k(n-k)}\right)$$
$$= O\left(\sum_{0 < k < n} \frac{1}{n} \left(\frac{1}{k} + \frac{1}{n-k}\right) \log^2 n\right)$$
$$= O\left(H_{n-1} \frac{\log^2 n}{n}\right) = O\left(\frac{\log^3 n}{n}\right)$$

We can take out the big-O because all summands are nonnegative.

#### Problem

Find an asymptotic estimate for 
$$g_n = [z^n]G(z)$$
 where  $G(z) = e^{\sum_{k \geqslant 1} \frac{z^n}{k^2}}$  .

#### Step 3: Manipulate the estimate obtained

We now rewrite our original recurrence by "pulling out the largest part":

$$ng_n = \sum_{k=0}^{n-1} \frac{g_k}{n} + \sum_{k=0}^{n-1} g_k \left( \frac{1}{n-k} - \frac{1}{n} \right) = \frac{1}{n} \sum_{k \ge 0} g_k - \frac{1}{n} \sum_{k \ge n} g_k + \frac{1}{n} \sum_{k=0}^{n-1} \frac{kg_k}{n-k}.$$

We conclude:

$$g_n = \frac{e^{\pi^2/6}}{n^2} + O\left(\frac{1}{n^2}O\left(\frac{\log^2 n}{n}\right)\right) + O\left(\frac{1}{n^2}O\left(\frac{\log^3 n}{n}\right)\right)$$
$$= \frac{e^{\pi^2/6}}{n^2} + O\left(\left(\frac{\log n}{n}\right)^3\right).$$



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### 1 O Manipulation

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Suppose we must approximate a function of the form:

$$f(n)=\sum_k a_k(n).$$

- **1** For each *n*, split the *k*s between a dominant set  $D_n$  and a tail set  $T_n$ .
- 2 Find asymptotic estimates  $a_k(n) = b_k(n) + O(c_k(n))$  which hold for  $k \in D_n$ .
- 3 Make sure that the following three sums are all "small":

$$S_{a}(n) = \sum_{k \in T_{n}} a_{k}(n); \ S_{b}(n) = \sum_{k \in T_{n}} b_{k}(n); \ S_{c}(n) = \sum_{k \in D_{n}} |c_{k}(n)|.$$

Then the following estimate holds:

$$f(n) = \sum_{k} b_{k}(n) + O(S_{a}(n)) + O(S_{b}(n)) + O(S_{c}(n)).$$



## Trading tails: The rationale

Suppose we performed the three steps in the previous slide. Then:

For k in the dominant set, we have the estimate:

$$\sum_{k \in D_n} a_k(n) = \sum_{k \in D_n} (b_k(n) + O(c_k(n)))$$
$$= \left(\sum_{k \in D_n} b_k(n)\right) + O(S_c(n))$$

Note that, to pass the big-O outside the summation, we needed the absolute values of the  $c_k(n)$ .

For k in the tail set, we have the estimate:

$$\sum_{k \in T_n} a_k(n) = \sum_{k \in T_n} (b_k(n) + a_k(n) - b_k(n))$$
$$= \left(\sum_{k \in T_n} b_k(n)\right) + O(S_a(n)) + O(S_b(n))$$

We did not need the absolute values of  $a_k(n)$  and  $b_k(n)$ , because they did appear explicitly in the summation, not inside a big-O.



Find an asymptotic estimate for the following sequence:

$$L_n = \sum_{k \ge 0} \frac{\ln(n+2^k)}{k!}$$

We immediately observe that, if k is "small" and n is "large", then:

$$\ln(n+2^{k}) = \ln\left(n \cdot \left(1+\frac{2^{k}}{n}\right)\right) = \ln n + \frac{2^{k}}{n} - \frac{2^{2k}}{2n^{2}} + O\left(\frac{2^{3k}}{n^{3}}\right)$$

and we can surely use this approximation within the convergence radius of the Taylor series of  $\ln(1+z)$  at the origin of the complex plane. Such convergence radius is 1, so we require  $2^k < n$ , that is,  $k < \lfloor \lg n \rfloor$ .



Find an asymptotic estimate for the following sequence:

$$L_n = \sum_{k \ge 0} \frac{\ln(n+2^k)}{k!}$$

Step 1 For every  $n \ge 1$  we set:

 $D_n = [0: \lfloor \lg n \rfloor - 1]$  and  $T_n = \mathbb{N} \setminus D_n = \{\lfloor \lg n \rfloor, \lfloor \lg n \rfloor + 1, \ldots\}$ 

Step 2 For  $k \in D_n$  we write  $a_k(n) = b_k(n) + O(c_k(n))$  where:

$$\begin{aligned} a_k(n) &= \frac{\ln(n+2^k)}{k!}; \\ b_k(n) &= \frac{1}{k!} \left( \ln n + \frac{2^k}{n} - \frac{4^k}{2n^2} \right); \\ c_k(n) &= \frac{8^k}{n^3 \cdot k!}. \end{aligned}$$



Find an asymptotic estimate for the following sequence:

$$L_n = \sum_{k \ge 0} \frac{\ln(n+2^k)}{k!}$$

Step 3c The estimate for  $S_c(n)$  is immediate:

$$\sum_{0 \leqslant k < \lfloor \lg n \rfloor} \frac{8^k}{n^3 \cdot k!} \leqslant \frac{1}{n^3} \cdot \sum_{k \geqslant 0} \frac{8^k}{k!} = \frac{e^8}{n^3} = O\left(\frac{1}{n^3}\right).$$



Find an asymptotic estimate for the following sequence:

$$L_n = \sum_{k \ge 0} \frac{\ln(n+2^k)}{k!}$$

Step 3b To estimate  $S_b(n)$  we can replace  $\frac{2^k}{n} - \frac{4^k}{2n^2}$  with  $2^k + 4^k$  and find:

$$\left| \sum_{k \ge \lfloor \lg n \rfloor} b_k(n) \right| < \sum_{k \ge \lfloor \lg n \rfloor} \frac{\ln n + 2^k + 4^k}{k!} \\ < \frac{\ln n + 2^{\lfloor \lg n \rfloor} + 4^{\lfloor \lg n \rfloor}}{\lfloor \lg n \rfloor!} \cdot \sum_{k \ge 0} \frac{4^k}{k!}.$$

The last sum is  $e^4$ , while the fraction is  $O\left(\frac{n^2}{\lfloor \lg n \rfloor!}\right)$  because of the summand  $4^{\lfloor \lg n \rfloor} = n^2$ . But  $\lfloor \lg n \rfloor!$  grows faster than any power of n, so  $S_b(n)$  is really very small for large n.

Find an asymptotic estimate for the following sequence:

$$L_n = \sum_{k \ge 0} \frac{\ln(n+2^k)}{k!}$$

Step 3a The estimate for  $S_a(n)$  is easy to compute: for  $n \ge 2$  and  $k \ge 1$ ,

$$\sum_{k \ge \lfloor \lg n \rfloor} \frac{\ln(n+2^k)}{k!} \le \sum_{k \ge \lfloor \lg n \rfloor} \frac{\ln(n \cdot 2^k)}{k!} < \sum_{k \ge \lfloor \lg n \rfloor} \frac{k+\ln n}{k!} = O\left(\frac{1}{\lfloor \lg n \rfloor!}\right)$$

because  $\ln 2^k < \ln e^k = k$ .



Find an asymptotic estimate for the following sequence:

$$L_n = \sum_{k \ge 0} \frac{\ln(n+2^k)}{k!}$$

We can now summarize:

$$L_n = \sum_{k \ge 0} \frac{1}{k!} \left( \ln n + \frac{2^k}{n} - \frac{4^k}{2n^2} \right) + O\left(\frac{1}{\lfloor \lg n \rfloor \rfloor}\right) + O\left(\frac{n^2}{\lfloor \lg n \rfloor \rfloor}\right) + O\left(\frac{1}{n^3}\right)$$
  
=  $e \ln n + \frac{e^2}{n} - \frac{e^4}{2n^2} + O\left(\frac{1}{n^3}\right).$ 



## Next section

### 1 O Manipulation

Two Asymptotic Tricks
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## Euler's summation formula

#### Theorem

If f(x) is differentiable m times in an open interval which contains [a:b], then:

$$\sum_{k \leq k < b} f(k) = \int_{a}^{b} f(x) \, dx + \sum_{k=1}^{m} \frac{B_{k}}{k!} \left( f^{(k-1)}(b) - f^{(k-1)}(a) \right) + R_{m},$$

where  $B_k$  is the kth Bernoulli number and where:

$$R_m = (-1)^{m+1} \int_a^b \frac{B_m(\{x\})}{m!} f^{(m)}(x) \, dx,$$

where, in turn,  $B_m(x) = \sum_k {m \choose k} B_k x^{m-k}$  is the *m*th Bernoulli polynomial.



## Euler's summation formula: A special case

If  $f(x) = x^{m-1}$  with m positive integer, then Euler's summation formula becomes:

$$\sum_{a \leqslant k < b} x^{m-1} = \frac{b^m - a^m}{m} + \sum_{k=1}^m \frac{B_k}{k!} (m-1)^{\underline{k-1}} (b^{m-k} - a^{m-k}) + 0$$
$$= \frac{1}{m} \sum_{k=0}^m B_k \frac{m(m-1)^{\underline{k-1}}}{k!} (b^{m-k} - a^{m-k}) \text{ because } B_0 = 1$$
$$= \frac{1}{m} \sum_{k=0}^m \binom{m}{k} (b^{m-k} - a^{m-k})$$



### Euler's summation formula: The rationale

Let  $\sum$ ,  $\Delta$ ,  $\int$  and D be the operators of summation, difference, integration, and differentiation, respectively.

- Suppose that *f* is smooth: that is, it has derivatives of any order.
- Taylor's formula tells us that:  $f(x + \varepsilon) = f(x) + f'(x)\varepsilon + \frac{f''(x)}{2}\varepsilon^2 + ...$
- For  $\varepsilon = 1$ , and writing  $D^k f$  in place of  $f^{(k)}$ , this becomes:

$$\Delta f(x) = \sum_{k \ge 1} \frac{D^k f(x)}{k!} = (e^D - 1)f(x),$$

with what looks like a little abuse of notation ....

• Now, if  $\Delta = e^D - 1$ , then its inverse  $\Sigma$  must be:

$$\sum = \frac{1}{e^D - 1} = \frac{1}{D} \cdot \frac{D}{e^D - 1} = \frac{1}{D} \cdot \left( 1 + \sum_{k \ge 1} \frac{B_k}{k!} D^k \right) = \int + \sum_{k \ge 1} \frac{B_k}{k!} D^{k-1},$$

which is Euler's summation formula with an infinite sum and no remainder.

TAL TECH We give the proof by induction on  $m \ge 1$  with a = 0 and b = 1. We only need this case, because if  $a \le c \le b$ , then:

 $\sum_{a\leqslant k< b} f(k) = \sum_{a\leqslant k< c} f(k) + \sum_{c\leqslant k< b} f(k);$ 

$$\int_{a}^{b} f(x) \, dx = \int_{a}^{c} f(x) \, dx + \int_{c}^{b} f(x) \, dx, \text{ and similar for } \int_{a}^{b} \frac{B_{m}(\{x\})}{m!} f^{(m)}(x) \, dx;$$

•  $f^{(k-1)}(b) - f^{(k-1)}(a) = (f^{(k-1)}(c) - f^{(k-1)}(a)) + (f^{(k-1)}(b) - f^{(k-1)}(c))$  for every k from 1 to m; and

• if 
$$a \neq 0$$
 we can replace  $f(x)$  with  $g(x) = f(x+a)$ .

So what we need to prove is:

$$f(0) = \int_0^1 f(x) \, dx + \sum_{k=1}^m \frac{B_k}{k!} \left( f^{(k-1)}(1) - f^{(k-1)}(0) \right) - (-1)^m \int_0^1 \frac{B_m(x)}{m!} f^{(m)}(x) \, dx$$

for every function f differentiable m times in (s, t) for some s < 0 and t > 1.



We give the proof by induction on  $m \ge 1$  with a = 0 and b = 1.

Base case: m = 1. Then  $B_1 = -1/2$ ,  $B_1(x) = x - 1/2$ , and the formula becomes:

$$f(0) = \int_0^1 f(x) \, dx - \frac{1}{2} (f(1) - f(0)) + \int_0^1 \left( x - \frac{1}{2} \right) f'(x) \, dx \, dx$$

or equivalently,

$$\frac{f(0)+f(1)}{2} = \int_0^1 f(x) \, dx + \int_0^1 \left(x - \frac{1}{2}\right) f'(x) \, dx \, ,$$

But the right-hand side is precisely:

$$\int_0^1 (f(x) + xf'(x)) \, dx - \frac{1}{2} \int_0^1 f'(x) \, dx = xf(x)|_0^1 - \frac{1}{2}(f(1) - f(0))$$
  
=  $f(1) - \frac{1}{2}f(1) + \frac{1}{2}f(0)$   
=  $\frac{f(0) + f(1)}{2}$ .



We give the proof by induction on  $m \ge 1$  with a = 0 and b = 1.

Induction: Suppose the thesis is true for  $m-1 \ge 1$ . Proving it for m is equivalent to proving that  $R_m = R_{m-1} - \frac{B_m}{m!} (f^{(m-1)}(1) - f^{(m-1)}(0))$ , that is:

$$(-1)^{m}B_{m}\left(f^{(m-1)}(1)-f^{(m-1)}(0)\right) = m\int_{0}^{1}B_{m-1}(x)f^{(m-1)}(x)\,dx$$
$$+\int_{0}^{1}B_{m}(x)f^{(m)}(x)\,dx$$



We give the proof by induction on  $m \ge 1$  with a = 0 and b = 1.

Induction: Suppose the thesis is true for  $m-1 \ge 1$ . Proving it for m is equivalent to proving that  $R_m = R_{m-1} - \frac{B_m}{m!} (f^{(m-1)}(1) - f^{(m-1)}(0))$ , that is:

$$(-1)^{m}B_{m}\left(f^{(m-1)}(1)-f^{(m-1)}(0)\right) = m\int_{0}^{1}B_{m-1}(x)f^{(m-1)}(x)\,dx$$
$$+\int_{0}^{1}B_{m}(x)f^{(m)}(x)\,dx$$

Now, if  $B_m(x) = \sum_k {m \choose k} B_k x^{m-k}$  (which is the case) then:

$$\frac{d}{x}B_{m}(x) = \sum_{k} {\binom{m}{k}(m-k)B_{k}x^{m-1-k}} \\
= \sum_{k} \frac{m^{\underline{k}}(m-k)}{k!}B_{k}x^{m-1-k} \\
= \sum_{k} \frac{m(m-1)^{\underline{k}}}{k!}B_{k}x^{m-1-k} \\
= m\sum_{k} {\binom{m-1}{k}}B_{k}x^{m-1-k} = mB_{m-1}(x)$$



В

We give the proof by induction on  $m \ge 1$  with a = 0 and b = 1.

Induction: Suppose the thesis is true for  $m-1 \ge 1$ . Proving it for m is equivalent to proving that  $R_m = R_{m-1} - \frac{B_m}{m!}(f^{(m-1)}(1) - f^{(m-1)}(0))$ , that is:

$$(-1)^{m}B_{m}\left(f^{(m-1)}(1)-f^{(m-1)}(0)\right) = m\int_{0}^{1}B_{m-1}(x)f^{(m-1)}(x)\,dx$$
$$+\int_{0}^{1}B_{m}(x)f^{(m)}(x)\,dx$$

ut since 
$$\frac{d}{dx}B_m(x) = mB_{m-1}(x)$$
, the right-hand side is:  

$$\int_0^1 \left( \left( \frac{d}{dx}B_m(x) \right) f^{(m-1)}(x) + B_m(x) \frac{d}{dx} f^{(m-1)}(x) \right) dx$$

$$= \int_0^1 \left( \frac{d}{dx}B_m(x) f^{(m-1)}(x) \right) dx$$

$$= B_m(1) f^{(m-1)}(1) - B_m(0) f^{(m-1)}(0)$$



We give the proof by induction on  $m \ge 1$  with a = 0 and b = 1.

Induction: Suppose the thesis is true for  $m-1 \ge 1$ . Proving it for m is equivalent to proving that  $R_m = R_{m-1} - \frac{B_m}{m!} (f^{(m-1)}(1) - f^{(m-1)}(0))$ , that is:

$$(-1)^{m}B_{m}\left(f^{(m-1)}(1)-f^{(m-1)}(0)\right)=B_{m}(1)f^{(m-1)}(1)-B_{m}(0)f^{(m-1)}(0)$$

But the above can be rewritten:

$$g(1)f^{(m-1)}(1) - g(0)f^{(m-1)}(0) = 0$$
 with  $g(x) = B_m(x) - (-1)^m B_m$ 

and this must hold for every f differentiable m times: the only possibility is that g(0) = g(1) = 0, that is,

$$B_m(0) = B_m(1) = (-1)^m B_m$$

and this must hold whatever  $m \ge 2$  is.



We give the proof by induction on  $m \ge 1$  with a = 0 and b = 1.

Induction: Suppose the thesis is true for  $m-1 \ge 1$ . Proving it for m is equivalent to proving that

$$B_m(0) = B_m(1) = (-1)^m B_m, \ m \ge 2.$$

But  $B_m(0) = B_m(1)$  for  $m \ge 2$  follows directly from the defining equation of Bernoulli numbers:

$$\sum_{k} \binom{m}{k} B_{k} = B_{m} + [m = 1] \text{ for every } m \ge 1$$

and the  $(-1)^m$  sign is not a problem, because for odd m > 1 it is  $B_m = 0$ . Q.E.D.



As  $B'_m(x) = mB_{m-1}(x)$  for every  $m \ge 0$ , from our discussion follows that:

$$\int_0^1 B_m(x) \, dx = \frac{B_{m+1}(1) - B_{m+1}(0)}{(m+1)!} = 0 \quad \text{for every } m \ge 1.$$

Then the remainder  $R_m$  is the integral of the product of an *m*th derivative with a function of average zero, everything divided by a factorial:

$$R_m = \frac{(-1)^{m+1}}{m!} \int_a^b B_m(\{x\}) f^{(m)}(x) \, dx = \frac{(-1)^{m+1}}{m!} \sum_{a \le k < b} \int_0^1 B_m(x) f^{(m)}(x+k) \, dx$$

Such a quantity has good chances to be small, even if the  $B_m$  grow very large. Actually, as  $\sum_{m \ge 0} \frac{B_m}{m!} z^m = \frac{z}{e^z - 1}$  and the right-hand side is differentiable in the entire complex plane, the left-hand side has infinite convergence radius, so  $B_m/m!$  vanishes faster than exponentially.



We have observed that  $B'_m(x) = mB_{m-1}(x)$  for every  $m \ge 1$ . If  $x \in [0,1]$  then:

- **B**<sub>1</sub>(x) = x 1/2 is negative in (0,1/2) and positive in (1/2,1), and  $B_1(1-x) = -B_1(x)$ .
- Then  $B_2(x)$  is decreasing in (0, 1/2) and increasing in (1/2, 1), and by comparing derivatives,  $B_2(1-x) = B_2(x)$ . Also,  $B_2(0) = B_2 = 1/6 > 0$ .
- By comparing derivatives,  $B_3(1-x) = -B_3(x)$ . As  $B_3(0) = B_3 = 0$  and  $B'_3(x) = 3B_2(x) > 0$  near 0,  $B_3(x)$  is positive in (0,1/2) and negative in (1/2,1).
- Then  $B_4(x)$  is increasing in (0, 1/2) and decreasing in (1/2, 1), and by comparing derivatives,  $B_4(1-x) = B_4(x)$ . Also,  $B_4(0) = B_4 = -1/30 < 0$ .
- By comparing derivatives,  $B_5(1-x) = -B_5(x)$ . As  $B_5(0) = B_5 = 0$  and  $B'_5(x) = 5B_4(x) < 0$  near 0,  $B_5(x)$  is negative in (0, 1/2) and positive in (1/2, 1).
- And so on . . .



# Behavior of $B_m(x)$ at x = 1/2

From the previous slide we deduce that for  $x \in [0,1]$ ,  $|B_{2m}(x)|$  is maximum at either x = 0 or x = 1/2.

#### Lemma

For every  $m \ge 0$ ,  $B_m(1/2) = (2^{1-m}-1) B_m$ .



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#### Lemma

For every  $m \ge 0$ ,  $B_m(1/2) = (2^{1-m}-1)B_m$ .

Proof  $B_m(x) = \sum_k {m \choose k} B_k x^{m-k}$  is the *m*th term of the binomial convolution of  $\langle B_m \rangle$  and  $\langle x^m \rangle$ . Then:

$$\sum_{m\geq 0}\frac{B_m(x)}{m!}z^m=\frac{ze^{xz}}{e^z-1},$$

which for x = 1/2 becomes:

$$\sum_{n=0}^{\infty} \frac{B_m(1/2)}{m!} z^m = \frac{ze^{z/2}}{e^z - 1}$$

$$= 2\frac{(z/2)(e^{z/2} + 1)}{e^z - 1} - \frac{z}{e^z - 1}$$

$$= 2\sum_{m \ge 0} \frac{B_m}{m!} \left(\frac{z}{2}\right)^m - \sum_{m \ge 0} \frac{B_m}{m!} z^m$$

$$= \sum_{m \ge 0} \frac{(2^{1-m} - 1)B_m}{m!} z^m$$



# Behavior of $B_m(x)$ at x = 1/2

From the previous slide we deduce that for  $x \in [0,1]$ ,  $|B_{2m}(x)|$  is maximum at either x = 0 or x = 1/2.

#### Lemma

For every  $m \ge 0$ ,  $B_m(1/2) = (2^{1-m}-1)B_m$ .

As  $|2^{1-m}-1| < 1$ , we conclude:

#### Corollary

For every  $x \in [0,1]$  and integer  $m \geqslant 1$ ,  $|B_{2m}(x)| \leqslant |B_{2m}| = (-1)^{m-1}B_{2m}$ .



### Euler's summation formula and asymptotics: Estimates

Let us write Euler's summation formula again:

$$\sum_{a\leqslant k< b} f(k) = \int_a^b f(x) \, dx + \sum_{k=1}^m \frac{B_k}{k!} \left( f^{(k-1)}(b) - f^{(k-1)}(a) \right) + \frac{(-1)^{m+1}}{m!} \int_a^b B_m(\{x\}) f^{(m)}(x) \, dx.$$

Since  $B_m = 0$  for m > 1 odd and  $B_1 = -\frac{1}{2}$ , we can only consider m even and rewrite:

$$\sum_{a \leqslant k < b} f(k) = \int_{a}^{b} f(x) dx - \frac{f(b) - f(a)}{2} + \sum_{k=1}^{m} \frac{B_{2k}}{(2k)!} \left( f^{(2k-1)}(b) - f^{(2k-1)}(a) \right) - \frac{1}{(2m)!} \int_{a}^{b} B_{2m}(\{x\}) f^{(2m)}(x) dx.$$

But for  $x \in [0,1]$  it is  $|B_{2m}(x)| \leq |B_{2m}|$ , and as  $\frac{(2\pi)^{2m}}{2} \frac{|B_{2m}|}{(2m)!} = \sum_{k \geq 1} \frac{1}{k^{2m}}$  (cf. Ch. 6)

$$\sum_{a \leq k < b} f(k) = \int_{a}^{b} f(x) \, dx - \frac{f(b) - f(a)}{2} + \sum_{k=1}^{m} \frac{B_{2k}}{(2k)!} \left( f^{(2k-1)}(b) - f^{(2k-1)}(a) \right) \\ + O\left( (2\pi)^{-2m} \right) \int_{a}^{b} |f^{(2m)}(x)| \, dx \, .$$

If  $f^{(2m)}$  is nonnegative in [a, b], then:

$$|R_{2m}| \leq \frac{|B_{2m}|}{(2m)!} \int_{a}^{b} f^{(2m)}(x) \, dx = \frac{|B_{2m}|}{(2m)!} \left( f^{(2m-1)}(b) - f^{(2m-1)}(a) \right)$$

But as  $B_{2m+1} = 0$  for  $m \ge 1$ , it is  $R_{2m} = R_{2m+1}$ , so the first discarded term when we approximate to the 2*m*th order instead of the (2m+2)nd must be  $R_{2m} - R_{2m+2}$ .

#### Lemma

If  $f^{(2m+2)}(x) \ge 0$  for every  $x \in [a,b]$ , then  $(-1)^m R_{2m} \ge 0$ .

Proof for a = 0 and b = 1: (general case follows easily)

• 
$$R_{2m} = R_{2m+1} = \frac{1}{(2m+1)!} \int_0^1 B_{2m+1}(x) f^{(2m+1)}(x) dx.$$

- As  $f^{(2m+2)} \ge 0$ ,  $f^{(2m+1)}$  is nondecreasing, and since  $B_{2m+1}$  is symmetric around x = 1/2, the second half of the sinusoid counts more than the first.
- For m even,  $B_{2m+1}$  is negative in (0,1/2) and positive in (1/2,1); for m odd,  $B_{2m+1}$  is positive in (0,1/2) and negative in (1/2,1). The thesis follows.



## Next section

### 1 O Manipulation

- Z Two Asymptotic Tricks
  Trick 1: Bootstrapping
  Trick 2: Trading tails
- 3 Euler's Summation Formula

#### 4 Final Summations

A bell-shaped summandStirling's approximation



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# A bell-shaped approximation

#### Problem

Give an asymptotic approximation of:

$$\Theta_n = \sum_k e^{-k^2/n}$$



## A bell-shaped approximation

#### Problem

Give an asymptotic approximation of:

$$\Theta_n = \sum_k e^{-k^2/n}$$

Let us think first of the behavior of the summand:

- The function  $f(x) = e^{-x^2}$  is maximum at x = 0 with f(0) = 1; stays near 1 in [-1,1]; and becomes very small very quickly for  $x \to \pm \infty$ .
- Then  $f_n(x) = e^{-x^2/n}$  stays near 1 for  $|x| \leq \sqrt{n}$ , and vanishes quickly for  $x \to \pm \infty$ .
- We then expect  $\Theta_n$  to be of the form  $\Theta_n \approx C \cdot \sqrt{n}$ .



Since Euler's summation formula holds for every  $a \le b$ , it also holds for the limits for  $a \to -\infty$  and  $b \to +\infty$  when they exist. This is the case for  $f_n(x) = e^{-x^2/n}$ , so:

$$\begin{split} \sum_{k} e^{-k^{2}/n} &= \int_{-\infty}^{+\infty} e^{-x^{2}/n} \, dx \\ &+ \sum_{k=1}^{m} \frac{B_{k}}{k!} \left( \lim_{x \to +\infty} f_{n}^{(k-1)}(x) - \lim_{x \to -\infty} f_{n}^{(k-1)}(x) \right) \\ &+ (-1)^{m+1} \int_{-\infty}^{+\infty} \frac{B_{m}(\{x\})}{m!} f_{n}^{(m)}(x) \, dx \, . \end{split}$$

Now, for  $f(x) = e^{-x^2}$  it is  $f^{(k)}(x) = P_k(x)e^{-x^2}$  for a polynomial  $P_m(x)$  of degree k. As  $f_n(x) = f(x/\sqrt{n})$ , we have  $f_n^{(k)}(x) = n^{-k/2}f^{(k)}(x/\sqrt{n}) \to 0$  for  $x \to \pm \infty$ , hence:

$$\sum_{k} e^{-k^{2}/n} = \sqrt{\pi n} + (-1)^{m+1} \int_{-\infty}^{+\infty} \frac{B_{m}(\{x\})}{m!} f_{n}^{(m)}(x) \, dx \text{ for suitable } C > 0 \, .$$



# The Gaussian integral

#### Theorem

$$\int_{-\infty}^{+\infty} e^{-x^2} dx = \sqrt{\pi}.$$



# The Gaussian integral

#### Theorem

$$\int_{-\infty}^{+\infty} e^{-x^2} \, dx = \sqrt{\pi} \, .$$

#### Proof:

$$\left(\int_{-\infty}^{+\infty} e^{-x^2} dx\right)^2 = \int_{-\infty}^{+\infty} e^{-(x^2+y^2)} dx dy$$
$$= \int_{\theta=0}^{2\pi} \int_{\rho=0}^{+\infty} e^{-\rho^2} \rho d\rho d\theta$$
$$= 2\pi \cdot \frac{1}{2} \int_{0}^{+\infty} e^{-t} dt$$
with the change of variable  $t = \rho^2$ 
$$= \pi \cdot \left[-e^{-t}\right]_{0}^{+\infty} = \pi.$$



# The Gaussian integral

#### Theorem

$$\int_{-\infty}^{+\infty} e^{-x^2} \, dx = \sqrt{\pi} \, .$$

#### Corollary

$$\int_{-\infty}^{+\infty} e^{-x^2/n} \, dx = \sqrt{\pi n} \, .$$



The absolute error we make by approximating  $\Theta_n$  with  $\sqrt{\pi n}$  is:

$$(-1)^{m+1} \int_{-\infty}^{+\infty} \frac{B_m(\{x\})}{m!} f_n^{(m)}(x) dx = \frac{(-1)^{m+1}}{n^{m/2}} \int_{-\infty}^{+\infty} \frac{B_m(\{x\})}{m!} f^{(m)}\left(\frac{x}{\sqrt{n}}\right) dx$$
  
=  $\frac{(-1)^{m+1}}{n^{(m-1)/2}} \int_{-\infty}^{+\infty} \frac{B_m(\{t\sqrt{n}\})}{m!} f^{(m)}(t) dt$   
with the change of variable  $t = x/\sqrt{n}$   
=  $O\left(n^{(1-m)/2}\right)$ 

because  $B_m(x)$  is bounded in [0,1] and  $\int_{-\infty}^{+\infty} |f^{(m)}(x)| dx$  is finite. But  $m \ge 1$  is arbitrary, because f(x) is smooth in  $\mathbb{R}$ : we conclude

$$\Theta_n = \sqrt{\pi n} + O\left(n^{-M}\right)$$
 for any  $M > 0$ .



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# Stirling's approximation

#### Theorem

$$\ln n! = \left(n + \frac{1}{2}\right) \ln n - n + \ln \sqrt{2\pi} + \frac{1}{12n} - \frac{1}{360n^3} + O\left(\frac{1}{n^5}\right)$$



# Stirling's approximation

#### Theorem

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Proof: (sketch; see the textbook for details)

1 Prove that there exists a constant  $\sigma$  such that:

$$\ln n! = \left(n + \frac{1}{2}\right) \ln n - n + \sigma + \frac{1}{12n} - \frac{1}{360n^3} + O\left(\frac{1}{n^5}\right).$$

2 Use the formula from point 1, the trading tails technique, and the approximation for  $\Theta_n$  from the previous section to prove that:

$$\sum_{k} \binom{2n}{k} = 2^{2n} \frac{\sqrt{2\pi}}{e^{\sigma}} \left( 1 + O\left(n^{-1/2 + 3\varepsilon}\right) \right) \text{ for } 0 < \varepsilon < \frac{1}{6}$$

