# ACHIEVING PERFORMANCE AND AVAILABILITY GUARANTEES WITH AMAZON SPOT INSTANCES 

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## INTRODUCTION (1)

- Cloud computing = pay-as-you go + "unlimited" capacity
- Amazon Elastic Cloud Compute (EC2)
- Web service providing resizable compute capacity
- Different instance types, e.g., small, large, xlarge servers
- Different purchasing options

1. On-Demand Instances

- Pay for resources by the hour
- No long-term commitments

2. Reserved Instances

- Upfront payment for reserving instances
- Discounted usage rate


## INTRODUCTION (2)

3. Spot Instances

- Users bid for unused resources at a "limit price" (the maximum price the bidder is willing to pay)
- Amazon gathers the bids and determines a clearing price ("spot price") based on the bids and the available capacity
- A bidder gets the required instances if his/her limit price is above the clearing price. In this case, the bidder pays the clearing price (not his/her limit price)
- The clearing price is updated as new bids arrive. If the clearing price goes above the bidder's limit price, the bidder's running spot instances are terminated.


## INTRODUCTION (3)

- In this presentation I will focus on "spot instances"
- They allow laaS providers to sell spare capacity (via auctions), thus improving resource utilization
- At the same time they enable laaS users to acquire resources at a price lower than ondemand price
- However, spot instances are terminated any time the "clearing price" goes above the user's "limit price"


## PROBLEM STATEMENT

"Can we use spot instances to run paid web services while achieving
performance and availability guarantees?"

## TALK OUTLINE



## PART 1: SPOT PRICE PREDICTION

- We must determine the optimal "limit price" that the SaaS provider should bid on the spot market
- The goal is to bid the lowest possible price capable of achieving the desired level of availability
- Let's see how some spot prices look like...


## SPOT PRICES (M1.SMALL)



## SPOT PRICES (M1.LARGE)



## SPOT PRICES (M1.XLARGE)



## $1^{\text {ST }}$ ATTEMPT

- Assume that prices follow patterns, e.g., day/night, week days/week ends, etc.
- Use triple exponential smoothing ("Holt-Winters") for predicting future prices



## WIKIPEDIA TRAFFIC



## WIKIPEDIA TRAFFIC: TIME SERIES DECOMPOSITION

The max irregular
component is $\sim 11 \%$ of the data

$\stackrel{3}{7}$

## PREDICTING WIKIPEDIA TRAFFIC WITH HOLT WINTERS

Predicting Wikipedia traffic with Holt-Winters


## PREDICTING SPOT PRICES WITH HOLT WINTERS (1)



## PREDICTING SPOT PRICES WITH HOLT WINTERS (2)


$\varphi$

## PREDICTING SPOT PRICES WITH HOLT WINTERS (3) <br> Days [5-9]



## XLARGE PRICES: TIME SERIES DECOMPOSITION

The max irregular component is $\sim 50 \%$ of the data


## $2^{\text {ND }}$ ATTEMPT (1)

- What about employing the distribution of the relative errors to correct the prediction?



## 2ND ATTEMPT (2)

CDF rel. error (Wikipedia)


## $2^{\text {ND }}$ ATTEMPT (3)

CDF rel. error. (xlarge)


## ONE MORE THING...

Dec. 25, 2010 - Feb. 23, 2011


Jun. 22, 2011 - Sept. 23, 2011


- ... they look quite different, don't they?


## 3RD ATTEMPT

1. Employ the autocorrelation function (ACF) to determine the similarity between prices as a function of the time difference between them (lag)

Value at time $t$



- The autocorrelation function (ACF) of the prices confirms the lack of any relationship between prices over the time

2. Use a normal approximation to model prices

- There is some evidence that prices in similar markets are approximately normally distributed
$2^{\text {ND }}$ ATTEMPT

- The autocorrelation function (ACF) of the prices confirms the lack of any relationship between prices over the time

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$2^{\text {ND }}$ ATTEMPT

- The autocorrelation function (ACF) of the prices confirms the lack of any relationship between prices over the time

2. Use a normal approximation to model prices

- There is some evidence that prices in similar markets are approximately normally distributed


## SPOT PRICE DISTRIBUTION

|  | Data | Normal Approx. |
| :---: | :---: | :---: |
| Minimum | 0.029 | 0.02098 |
| $1{ }^{\text {st }}$ Quartile | 0.029 | 0.02877 |
| Median | 0.031 | 0.03020 |
| $3^{\text {rd }}$ Quartile | 0.031 | 0.03163 |
| Maximum | 0.085 | 0.04010 |
| Mean (1 ${ }^{\text {st }}$ moment) | 0.0302 | 0.03025 |
| Variance (2 ${ }^{\text {nd }}$ moment) | $4.503941 \times 10^{-6}$ | $4.515323 \times 10^{-6}$ |
| Skewness (3 ${ }^{\text {rd }}$ moment) | $1.877202 \times 10$ | $1.659259 \times 10^{-2}$ |
| Kurtosis (4 $4^{\text {th }}$ moment) | $4.700735 \times 10^{2}$ | 3.004374 |

m1.small (Linux/Unix), us-east1 region. Dec. 25, 2010 - Feb. 23, 2011



## PRICE PREDICTION ALGORITHM

(1) Collect price statistics, and compute mean and variance
(2) Compute ACF of the prices for lag I, with I being the prediction horizon
(3) If $A C F(I)>0.4$,
a) Predict the future prices using linear regression, $y=a+b x$, otherwise
b) Use the quantile function (inverse CDF) of the Normal distribution to predict the future prices

- The inverse CDF returns the value of $x$ such as $P(X \leq x)=p$
(4) Use the max value returned by (3) as a "limit price"
(5) In case of out-of-bid events, increase the bid by 40\%


## PREDICTION ALGORITHM PERFORMANCE

| Prediction <br> (hours) | Target <br> availability | Achieved <br> availability | Avg. big (\$/h) |
| :--- | :--- | :--- | :--- |
| 6 | $90 \%$ | $99.673 \%$ | 0.03170 |
| 6 | $99 \%$ | $99.673 \%$ | 0.03270 |
| 6 | $99.999 \%$ | $99.673 \%$ | 0.03463 |
| 12 | $90 \%$ | $99.675 \%$ | 0.03207 |
| 12 | $99 \%$ | $99.675 \%$ | 0.03307 |
| 12 | $99.999 \%$ | $99.675 \%$ | 0.03499 |
| 24 | $90 \%$ | $99.679 \%$ | 0.03253 |
| 24 | $99 \%$ | $99.679 \%$ | 0.03350 |
| 24 | $99.999 \%$ | $99.679 \%$ | 0.03533 |

m1.small (Linux/Unix), us-east1 region. Dec. 25, 2010 - Feb. 23, 2011. The price of "on-demand" instances is $0.085 \$ / \mathrm{h}$

## PART 2: MODEL AND SLA (1)

The SaaS provider

1. Pays the IaaS for renting computing resources
2. Charges its customer for executing $\mathcal{W S}$ transactions
3. Pays penalties to its customer for failing to meet performance and availability objectives

## MODEL AND SLA (2)...

1. For each accepted and completed request, the client pays a charge of c \$
2. Performance: if the avg. resp. time, $\beta$, over an interval of length $\boldsymbol{t}$ exceeds the threshold $q$, then the provider pays a penalty of $r_{1} \$$ for each job executed in the interval
3. Availability: the provider should pay a penalty of $\mathbf{r}_{2} \$$ for each rejected job
4. Disaster: the provider is liable to pay a $\mathbf{r}_{3} \$$ for every accepted job that is lost due to resources becoming unavailable
5. The provider pays $r_{4} \$ / h$ to rent each server

- The provider tries to optimize the average net revenue earned per unit time


## OR

Prob. that the avg. resp. time exceeds
Prob. that spot instances are terminated the threshold

$$
R=\gamma\left[c-r_{1} P(\beta>q)\right]-r_{2}(\lambda-\gamma)-r_{3} P\left(l_{p}<S_{D}\right) L-r_{4} n
$$

Rate at which jobs are rejected
Avg. no. of jobs in the system

- $r_{1}=$ penalty for "performance"
- $r_{2}=$ penalty for "availability"
- $r_{3}=$ penalty for "disaster"
- $r_{4}=$ server rental cost
- $\quad Y=$ rate at which jobs are accepted into the system


## POLICIES

Resource allocation and admission control

- Admission control defined by means of a threshold, $K$

1. "Threshold" policy

- The simultaneous optimization of $K$ and $n$ is non trivial
- Treats the system as an $M / M / n / K$ queue
- Uses a "Hill Climbing" algorithm to find the "best" value of $n$ and $K$

2. "Heuristic": assumes that the response time is dominated by the service time (true in "large" systems)

- Treats the system as an $G / / G / n$ queue
- $K=\infty$, e.g., no job is rejected


## PERFORMANCE EVALUATION - SETTINGS

| b | 0.5 sec. | Average service time |
| :--- | :--- | :--- |
| q | 1 sec. | Performance threshold |
| c | $2.951 \times 10^{-5} \$$ | Charge per job |
| $\mathrm{r}_{1}$ | $1.5 \times \mathrm{c}$ | Penalty (performance) |
| $\mathrm{r}_{2}$ | $2 \times \mathrm{c}$ | Penalty (availability) |
| $\mathrm{r}_{3}$ | $3 \times \mathrm{c}$ | Penalty (disaster) |
| $\mathrm{r}_{4}$ | 0.085 \$/h | Rental cost ("on-demand" instances) |
| t | 1 min. | Interval length (SLA evaluation) |

## REVENUE AS FUNCTION OF K AND N



$$
p\left(I_{p}<c_{p}\right)=0
$$

## REVENUE AS FUNCTION OF K AND N



$$
p\left(I_{p}<c_{p}\right)=0.25
$$

## REVENUE AS FUNCTION OF K AND N



$$
p\left(I_{p}<c_{p}\right)=0.5
$$

## REVENUE AS FUNCTION OF $\mathbf{R}_{\mathbf{1}}, \mathbf{R}_{\mathbf{2}}$ AND K



The system is saturated

- Have to choose between a long queue and rejecting many jobs


## REVENUE AS FUNCTION OF THE ARRIVAL RATE



1. The revenue increases with the load
2. The use spot instances makes the system more profitable
3. The "Heuristic" policy performs almost as good as the "Threshold" algorithm

Each point represents a run lasting 28 days
$c a^{2}=1, c s^{2}=1, \rho=7.5, \ldots, 25, p\left(l_{p}<c_{p}\right)=0.001$

## CONCLUSIONS

- I have discussed an approach aiming at maximizing the net revenue earned by a SaaS using spot instances to provide a web service to paying customers

1. How much to bid for resources on the spot market?
2. How many servers should be allocated for a given time period, and how many jobs to accept?

- The number of running servers, as well as the maximum number of jobs admitted into the system, have a significant effect on the earned revenues
- The optimal queue length is highly dependent on the availability level (e.g., shorter queue when the likelihood premature termination is high)

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