

Control and data dependency- based automated analysis of security protocols for confidentiality

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(joint work with Peeter Laud)

Problem statement



- Given the program (performing computations and exchanging messages over public channel),
- Working with some secret data
- No active adversary should be able to learn anything about the secret data
- Automatically determine whether the protocol is secure or not.

Existing Solutions

- Program with operational semantics.
- Adversary running the program and observing its outputs.
- Secrecy definition – adversary cannot distinguish between two identical programs working with different secret data.
- Properties of cryptographic primitives enable modification of protocol text
- It is possible to achieve the point when “secret” value is not used in the protocol anymore.
- Bruno Blanchet technique – uses this approach.
- Syntax trees or flow graphs are not the best program representation for transformations.

Our Technique Outline

- **Program in WHILE-style language.**
- Operational semantics.
- Adversary running the program and observing its outputs.
- Secrecy definition – adversary cannot distinguish between two identical programs working with different secret data.
- **Program dependency graph**
- Graph semantics – functional dependency of outputs on the inputs.
- Adversary supplying the inputs and observing the outputs.
- Secrecy definition – no functional dependency of the outputs on the secret data.

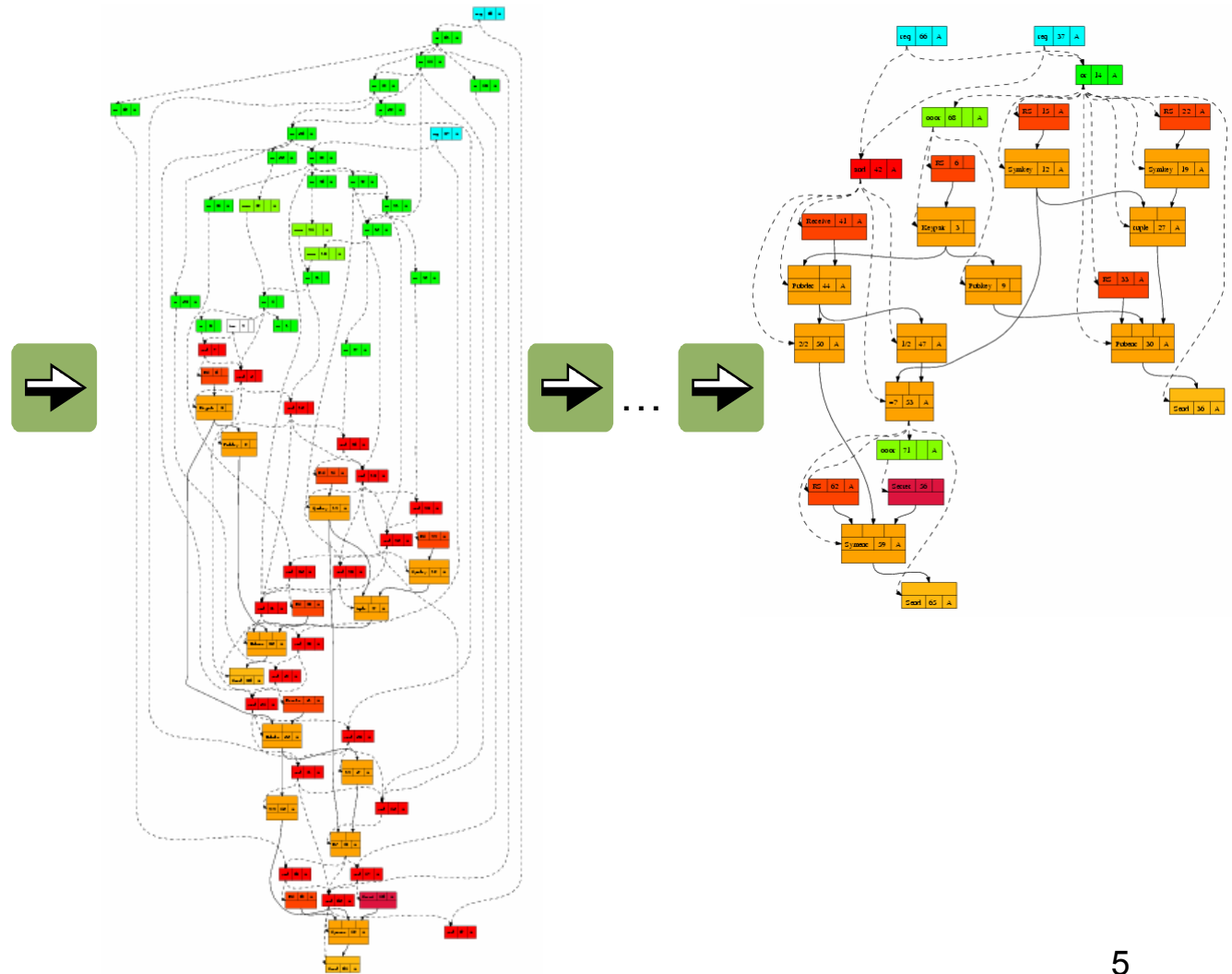
Technique Outline 2

```

kp:=gen_key_pair
pk:=get_pub_key(kp)

Replicate
(
  sk_1:=gen_s_key
  sk_2:=gen_s_key
  sks:=(sk_1,sk_2)
  sks_e:=aenc(pk,sks)
  send(sks_e)
  rsks_e:=receive
  rsks:=adecd(kp,rsks_e)
  rsk_1:=proj_1^2(rsks)
  rsk_2:=proj_2^2(rsks)
  check(rsk_1=sk_1)
  M_e:=senc(sk_2,M)
  send(M_e)
)

```



Program language semantics

- Structural operational semantics
- Program execution is set of transitions on the configuration set
- Configuration captures computation state at a given moment

The configuration is a tuple $\langle S, s, out \rangle$, where

- $S \in Stm \cup \{\epsilon\}$: unexecuted yet statements,
- $s \in State$: current state ($Var \rightarrow Val$),
- $out \in Val \cup \{ok\} \cup \{\perp\}$: output to the adversary

Program language semantics 2

$$\frac{\mathbb{A}[[a]]s \neq \perp}{\langle P\text{Asgn}(x, a), s, - \rangle \Rightarrow \langle \epsilon, s[x \mapsto \mathbb{A}[[a]]s], ok \rangle} (\text{Asgn}^{\text{ok}})$$

$$\frac{\mathbb{A}[[a]]s = \perp}{\langle P\text{Asgn}(x, a), s, - \rangle \Rightarrow \langle P\text{Stopped}, s, stuck \rangle} (\text{Asgn}^e)$$

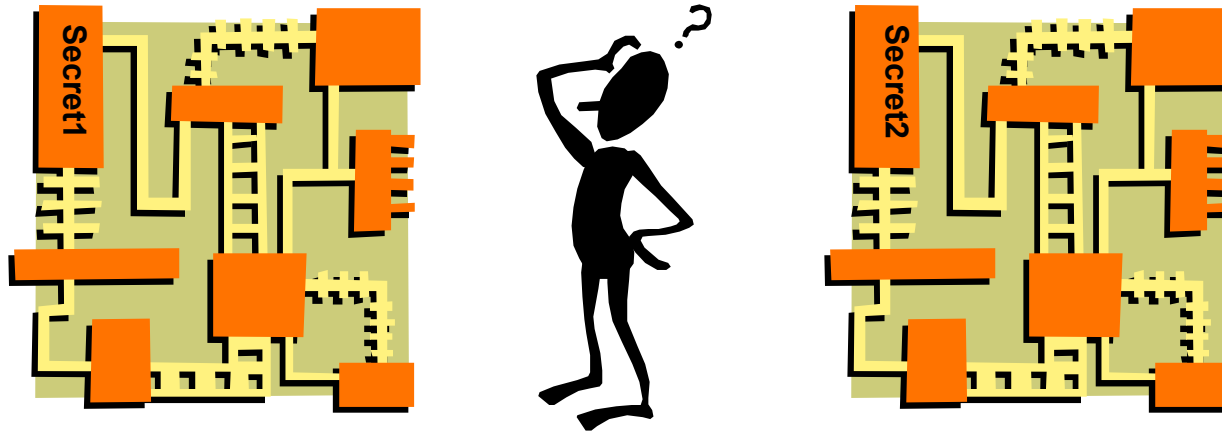
$$\frac{\langle S_1, s, - \rangle \Rightarrow \langle S'_1, s', out \rangle}{\langle P\text{Paral}(S_1, S_2), s, - \rangle \Rightarrow \langle P\text{Paral}(S'_1, S_2), s', out \rangle} (\text{Paral}^1)$$

$$\frac{\langle S_1, s, - \rangle \Rightarrow \langle \epsilon, s', out \rangle}{\langle P\text{Paral}(S_1, S_2), s, - \rangle \Rightarrow \langle S_2, s', out \rangle} (\text{Paral}^2)$$

$$\frac{\langle S_2, s, - \rangle \Rightarrow \langle S'_2, s', out \rangle}{\langle P\text{Paral}(S_1, S_2), s, - \rangle \Rightarrow \langle P\text{Paral}(S_1, S'_2), s', out \rangle} (\text{Paral}^3)$$

$$\frac{\langle S_2, s, - \rangle \Rightarrow \langle \epsilon, s', out \rangle}{\langle P\text{Paral}(S_1, S_2), s, - \rangle \Rightarrow \langle S_1, s', out \rangle} (\text{Paral}^4)$$

Security definition

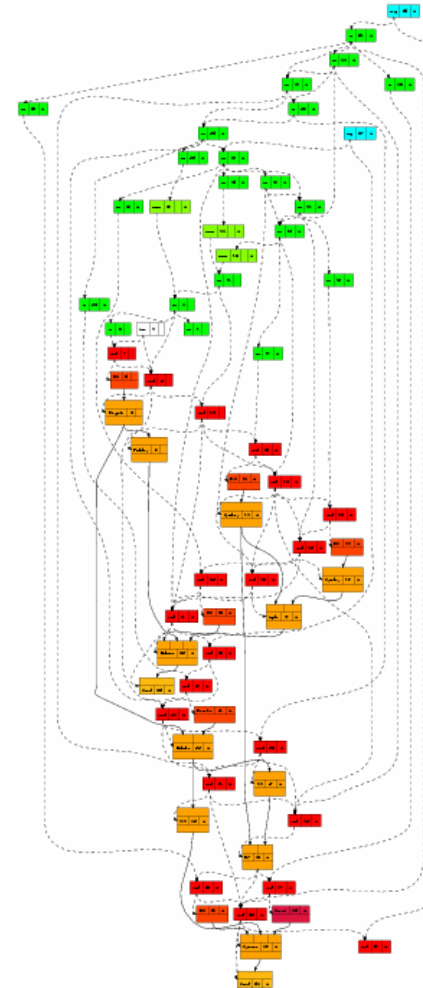


The protocol is considered secure if there's no non-negligible correlation between the final state of the adversary and the secret message

Translation to graphs

```
kp:=gen_key_pair
pk:=get_pub_key(kp)

Replicate
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  sk_2:=gen_s_key
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  sks_e:=aenc(pk,sks)
  send(sks_e)
  rsk_e:=receive
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  send(M_e)
)
```



Adversary “playing” with graph should observe at least same set of info as running the program.

Graph execution rules

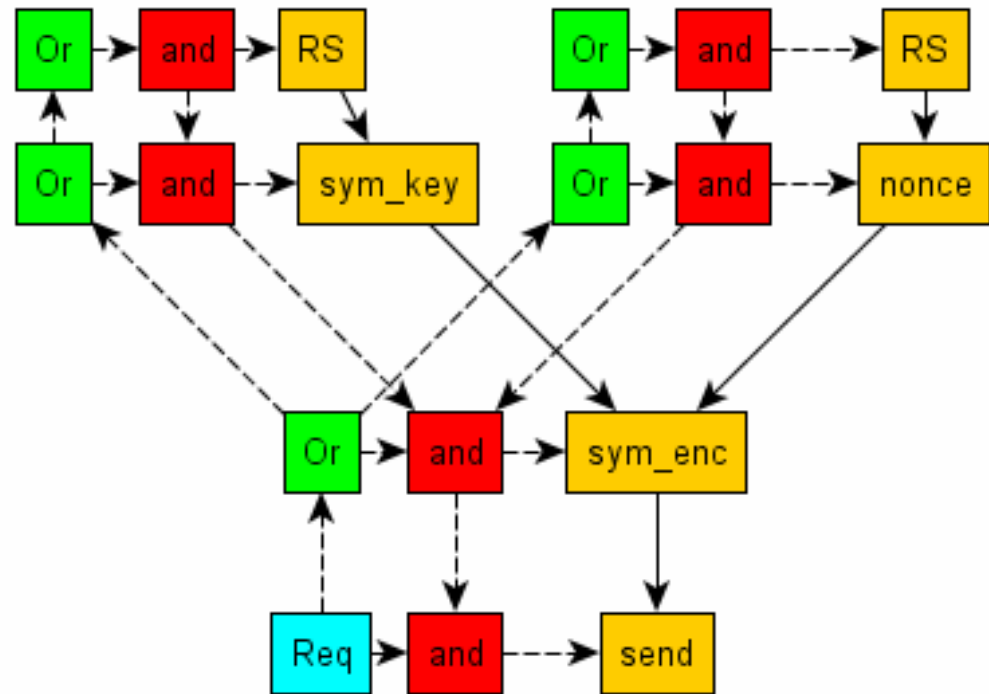
- Graph represents:
 - Operations occurring in the protocol (nodes)
 - Data dependencies (edges)
 - Control dependencies (edges)
 - Adversary interface (special type of nodes - Req)
 - Note: If replication is present, the Graph is infinite (each replicated operation is present in N copies).

Graph execution rules 2

- Execution rules
 - Adversary defines which outputs he'd like to see (sets true/false values to Req nodes), and which values are supplied to the inputs (values program gets from communication channels)
 - Now find the values of the outputs (using the graph semantics).

Graph execution rules - example

```
x := gen_sym_key  
y := random  
z := sym_enc (y, x)  
send (z)
```



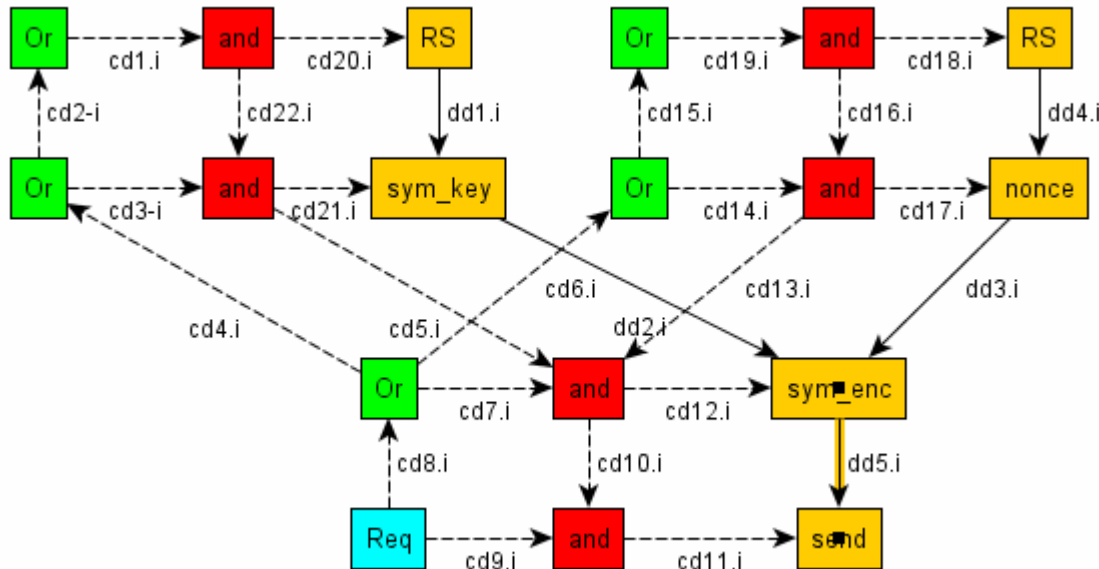
Graph semantics

- Edges (dependencies)
- Nodes (operations)
- Set of equations of type:

$$\text{Dep}_i = \text{Operation}(\text{Dep}_1, \dots, \text{Dep}_n)$$

 - One equation per node
 - Specifies how the dependencies are related – each dependency is a function of other dependencies
- Smallest solution of the set of equations – final values, subset of them visible to Adversary.
 - All Operations are monotone – so the smallest solution should exist.
- Security criteria – no functional dependency of outputs on secret.

Equation examples

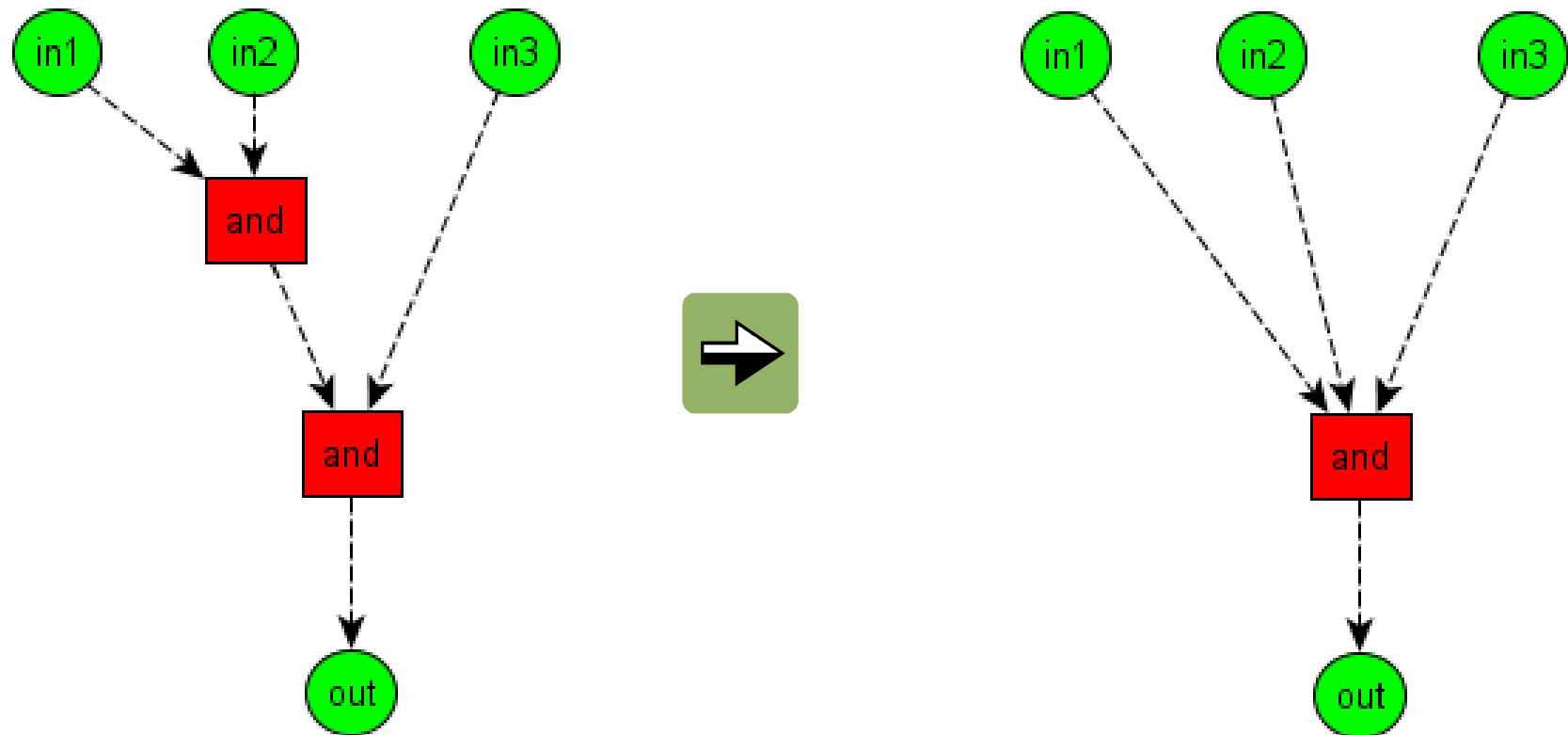


- Equation examples:
 - $dd3.i = \text{Nonce}(cd17.i, dd4.i)$
 - $cd17.i = \text{And}(cd14.i, cd16.i)$
 - $cd8.i = \text{Req} \leftarrow$ This value is supplied by Adversary.

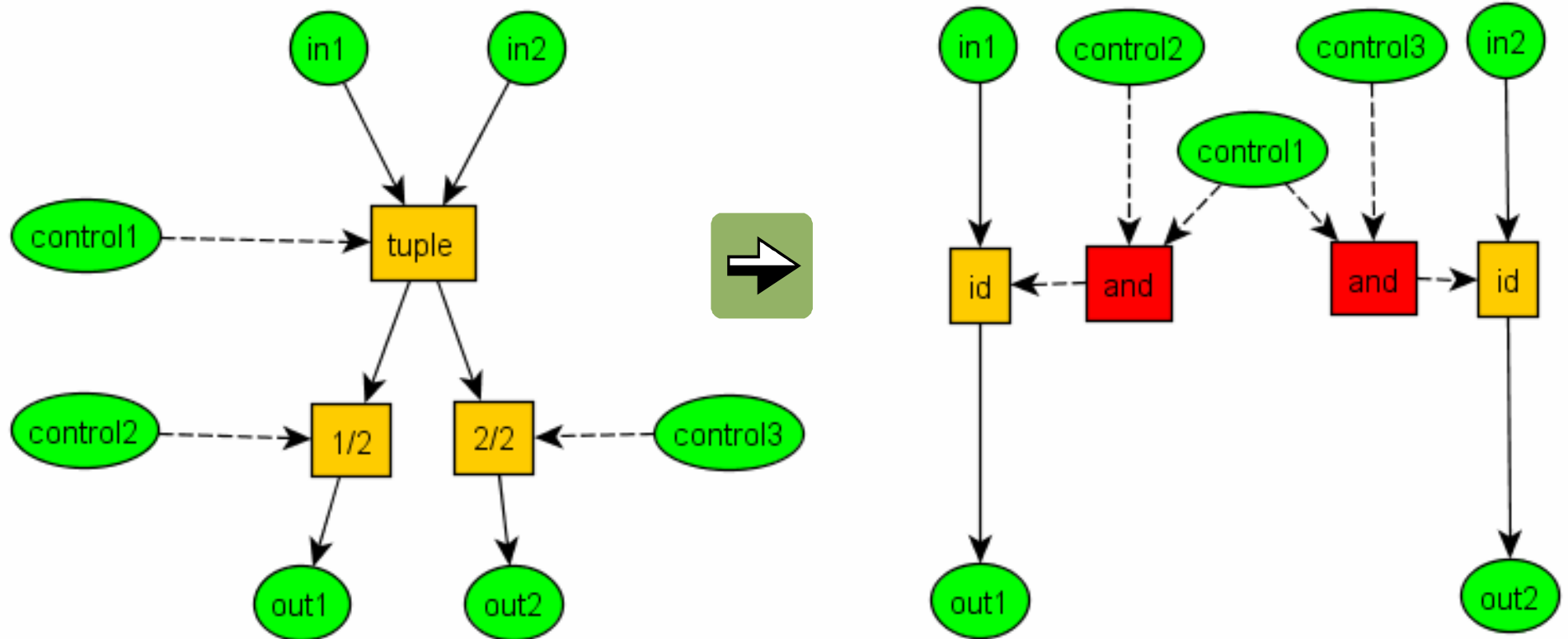
Graph transformation

- Transforming the graph / equations
- The smallest solution should stay indistinguishable from the original one.
- Transformations are based on the properties of the operation (including cryptographic primitives)

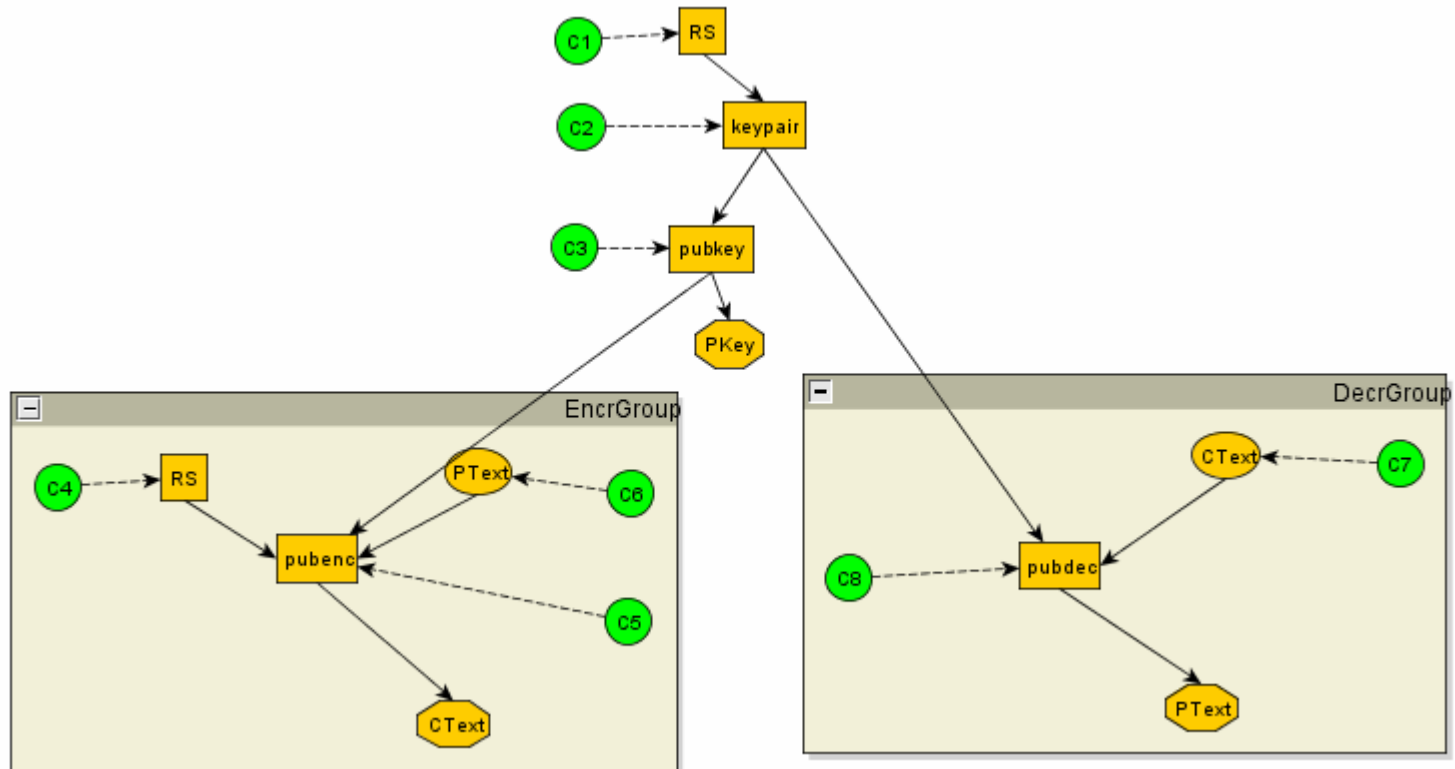
Control flow transformation example



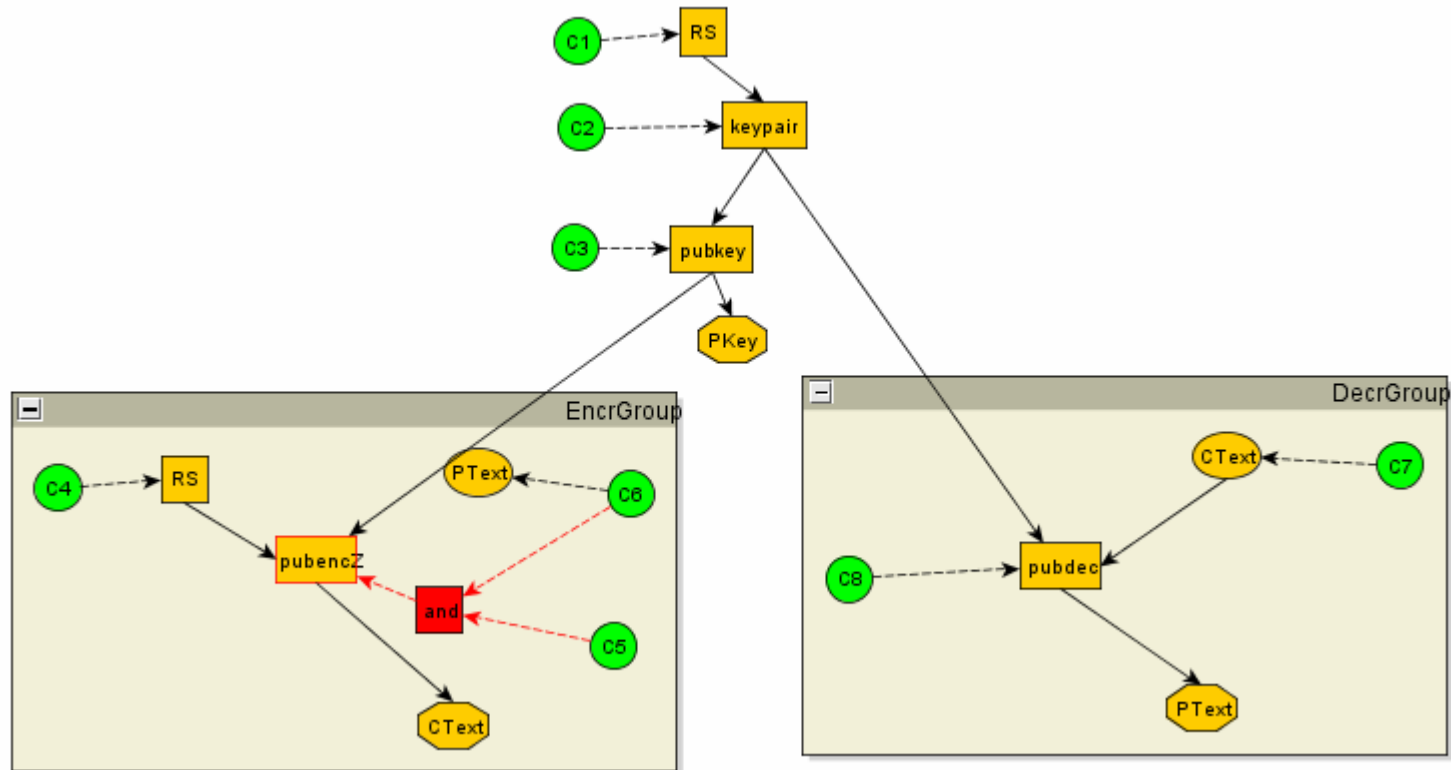
Operation transformation example



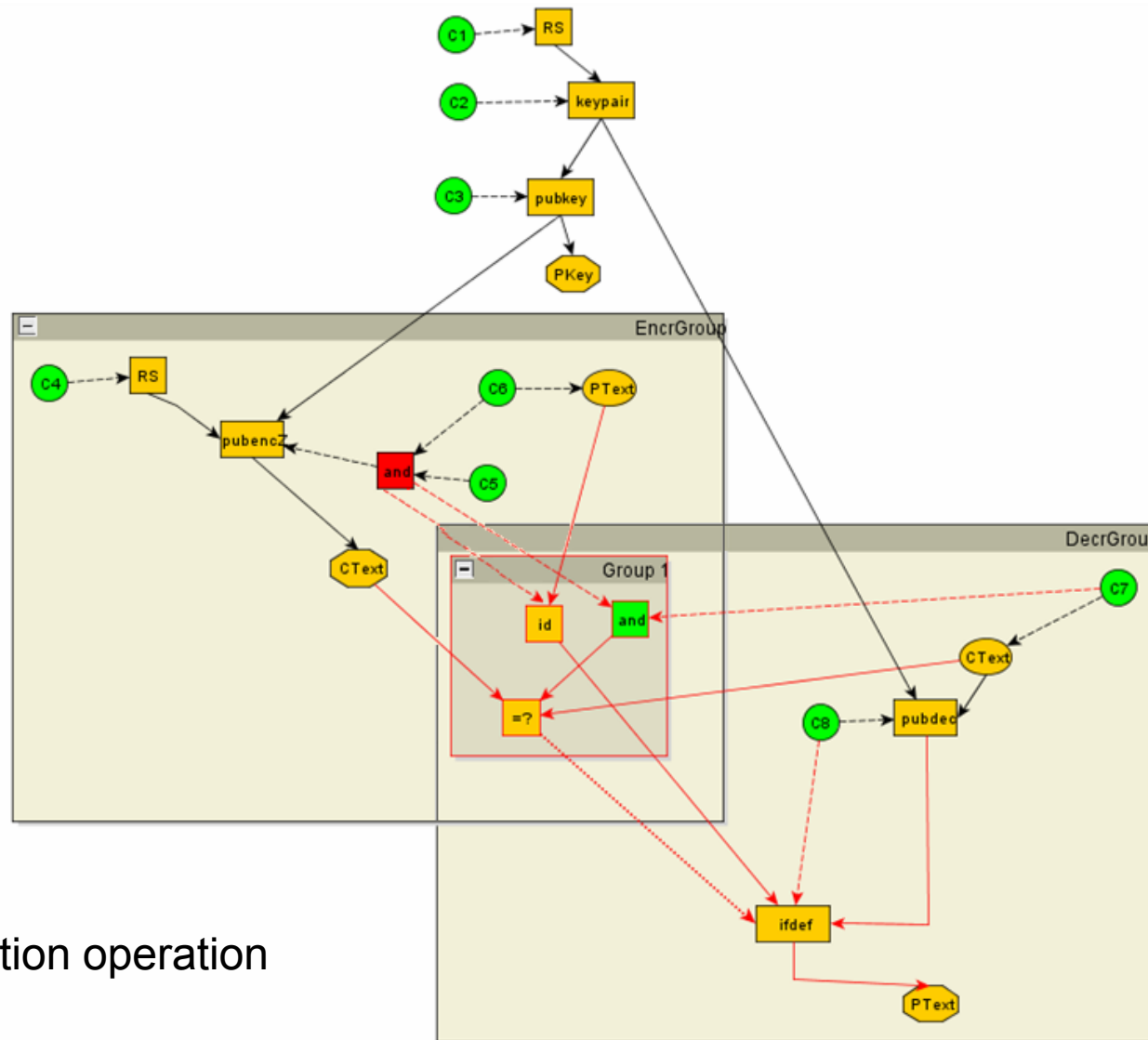
Pubenc transformation example - 1



Pubenc transformation example - 2

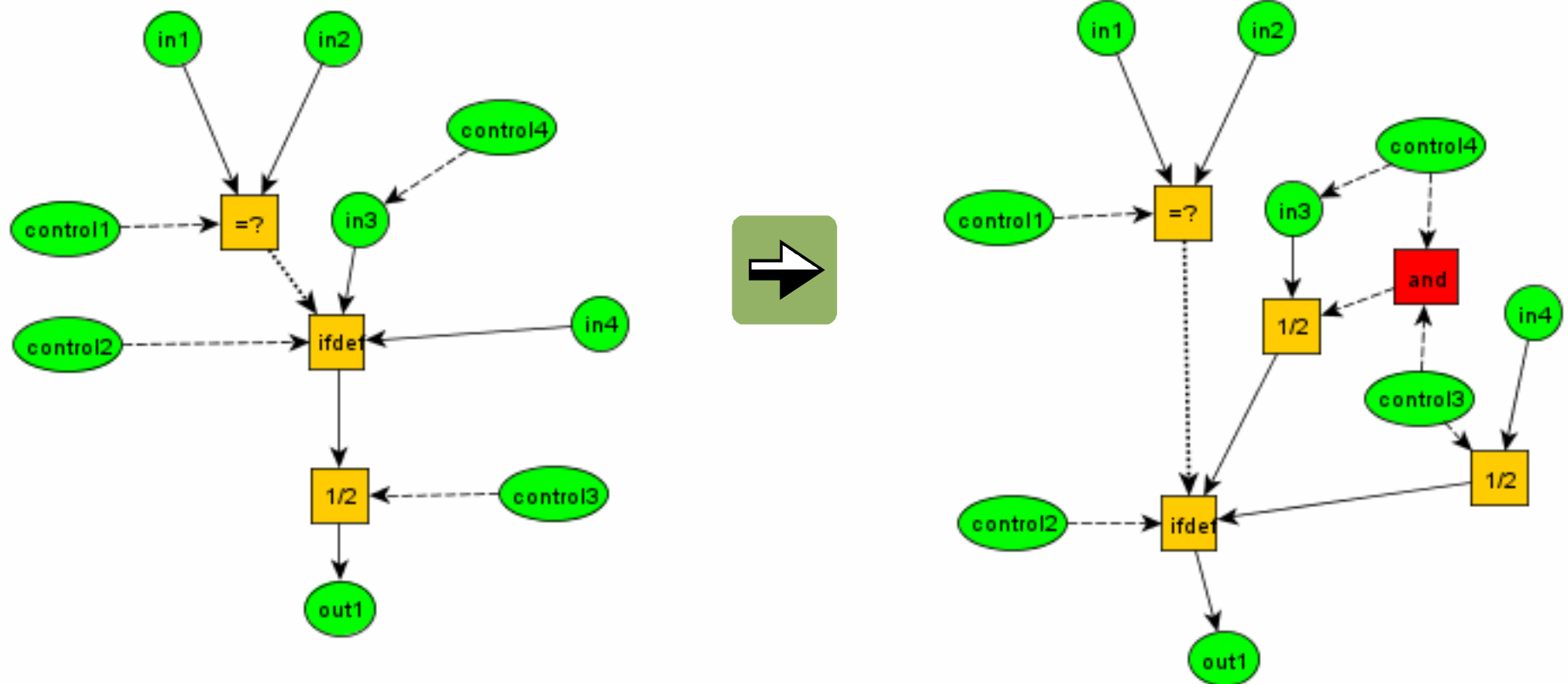


Pubenc transformation example - 3

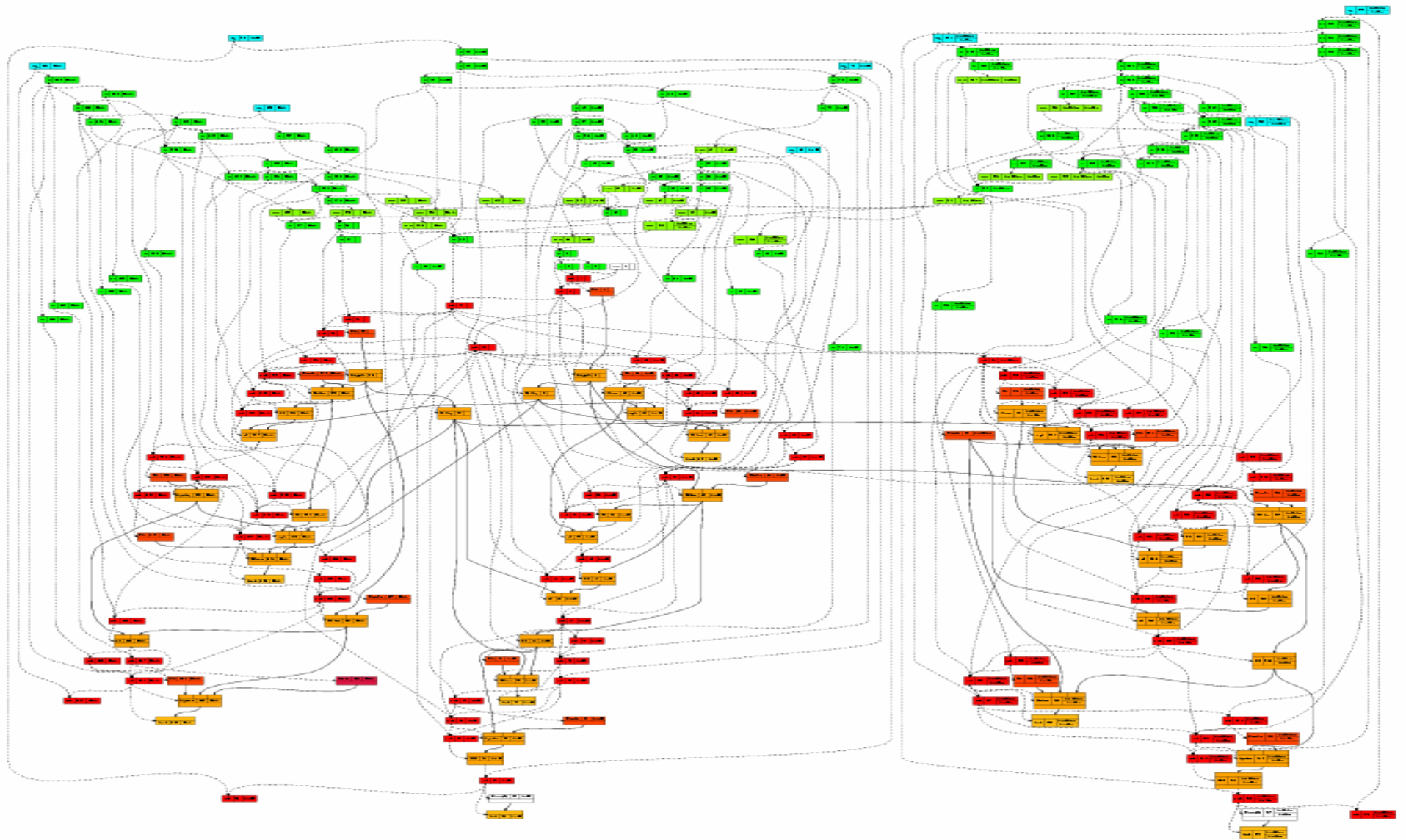


Ifdef – selection operation

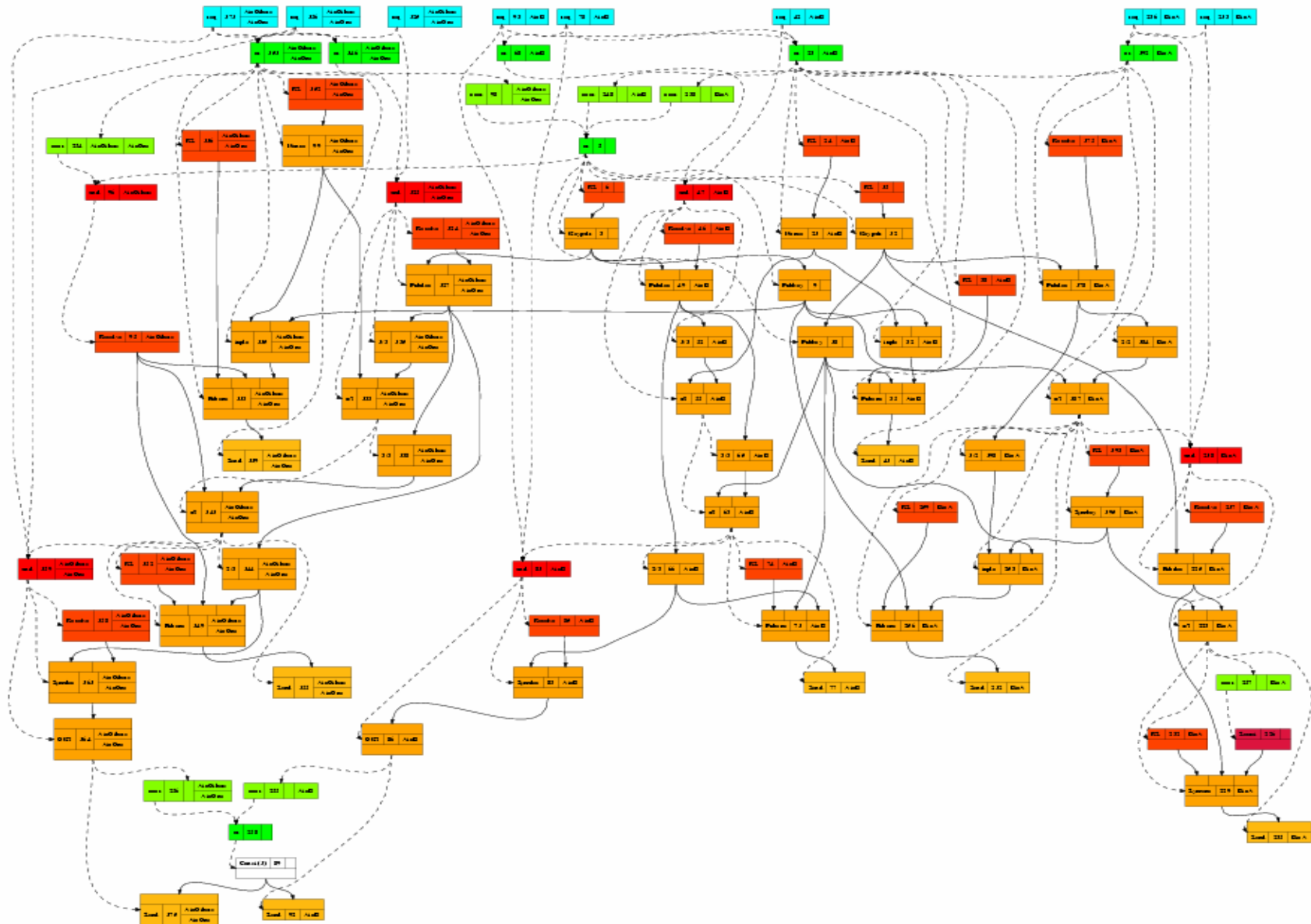
Operation transformation example



“Real” graph – a bit more complex 😊



“Real” graph – after some transforms applied – still complex.



Summary – what's done, what's left...

- Ready
 - A nice idea
 - Conceptual framework
 - Programming language semantics
 - Part of the graph semantics
 - Analyzer prototype
- To be done
 - Program \rightarrow graph translation correctness proof
 - Graph transformation correctness proof
 - Complete graph semantics
 - Fully functional analyzer

Thank you