Towards refined notions of computation: the global state example

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Overview

- Moggi's monadic approach to computational effects
- Lawvere theories
  and the computational effects they identify
- Refinement types
  and adding more detailed specifications
- Refinement types + Lawvere theories = ?
  on an example of refined global state

Moggi's monadic approach

Moggi's monadic approach

- Semantics of pure simply-typed lambda calculus:
  - take a cartesian-closed category $\mathcal{C}$
  - interpret base types $\alpha, \beta, \ldots$ as objects $[\![\alpha]\!], [\![\beta]\!], \ldots$
  - interpret product type as finite product structure on $\mathcal{C}$
  - interpret (pure) function type $\sigma \to \tau$ as the exponential $[\![\sigma]\!] \Rightarrow [\![\tau]\!]$
  - interpret value terms $\Gamma \vdash t : \sigma$ as morphisms $[\![\Gamma]\!] \longrightarrow [\![\sigma]\!]$

- Moggi's insight for impure languages:
  - use a strong monad $T : \mathcal{C} \longrightarrow \mathcal{C}$ to model computational effects
  - $T[\![\sigma]\!]$ stands for computations returning values from $[\![\sigma]\!]$
  - interpret impure function type $\sigma \rightharpoonup \tau$ as the Kleisli exponential $[\![\sigma]\!] \Rightarrow T[\![\tau]\!]$
  - interpret computations as Kleisli maps $[\![\Gamma]\!] \longrightarrow T[\![\sigma]\!]$

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Moggi's monadic approach

- Example monads proposed by Moggi
  - exceptions - $TX = X + E$
  - global state - $TX = (S \times X)^S$
    - *(stateful computations* $S \times X \longrightarrow S \times Y$*)*
  - local state - $(TX)_n = (\int^{m \in (n/I)} (S_m \times X_m))^{S_n}$
  - finite nondeterminism - $TX = \mathcal{F}^+ X$
  - continuations - $TX = R^{R^X}$

- Also possible to combine different monads, e.g.,
  - state plus exceptions - $TX = (S \times (X + E))^S$

Moggi's monadic approach

- Moggi's work gives us an elegant denotational semantics of computational effects
- However, this denotation does not tell us much about how to construct such effects
- We have to note their operational meaning and how such effects (e.g., state) are implemented in programming languages

Lawvere theories

Lawvere theories

- A countable Lawvere theory consists of:
  - a small category $\mathcal{L}$ with countable products
  - an id. on objects countable-product preserving functor

    $$J : \aleph_1^{op} \longrightarrow \mathcal{L}$$

  - *(where $\aleph_1$ is the skeleton of the category of countable sets)*

- Think of the hom $\mathcal{L}(n, 1)$ (abbrv. $\mathcal{L}(J(n), J(1))$) as a set of n-ary operations in the theory

- Then it suffices to give an algebraic theory as:
  - operations of are given by morphisms $op : O \longrightarrow I$
    - *(equivalently a family of operations $op_{i \in I} : O \longrightarrow 1$)*
  - equations are given by commuting diagrams

Models of Lawvere theories

- A model of a Lawvere theory $(\mathcal{L}, J)$ in a category $\mathcal{C}$ with countable products
  - is a countable product preserving functor $M : \mathcal{L} \longrightarrow \mathcal{C}$
- The models of $\mathcal{L}$ together with nat. transfs. :
  - form a category $Mod(\mathcal{L}, \mathcal{C})$ with $U : Mod(\mathcal{L}, \mathcal{C}) \longrightarrow \mathcal{C}$
  - having a left adjoint $F : \mathcal{C} \longrightarrow Mod(\mathcal{L}, \mathcal{C})$
  - the adjoint functors induce a monad $T = UF$
- For the purposes of this talk, we let $\mathcal{C} = \mathrm{Set}$
- To give a model M of $\mathcal{L}$ is equivalent to
  - giving a set $X = M1$
  - for every operation $op : O \longrightarrow I$ a morphism $X^O \longrightarrow X^I$
- Because
  - $M1$ determines MO up to coherent isomorphism
  - $MO \cong M(\prod_{o \in O} 1) \cong \prod_{o \in O}(M1) \cong (M1)^O$

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Global state example

- Plotkin and Power noticed that the global state monad is determined by the following countable Lawvere theory
- Countable set of values V and a finite set of locations Loc
- Take the set of states to be $S = V^{Loc}$
- The theory is freely generated by operations
  - $lookup : V \longrightarrow Loc$
  - $update : 1 \longrightarrow Loc \times V$

  subject to commuting diagrams expressed set-theoretically

  1. $lookup_{loc}(update_{loc,v}(x))_v = x$
  2. $lookup_{loc}(lookup_{loc}(x_{vv'})_v)_{v'} = lookup_{loc}(x_{vv})_v$
  3. $update_{loc,v}(update_{loc,v'}(x)) = update_{loc,v'}(x)$
  4. $update_{loc,v}(read_{loc}(x'_v)'_v) = update_{loc,v}(x_v)$
  5. $update_{loc,v}(update_{loc',v'}(x)) =$ $update_{loc',v'}(update_{loc,v}(x)) \quad (loc \neq loc')$
  6. ...

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  6. ...

Small detour into local state

- $(TX)_n = (\int^{m \in (n/Inj)} (S_m \times X_m))^{S_n}$
- Monad and algebra are given in category Set^{Inj}
    - *(Inj is the category of finite sets and injections)*

- $L_n = Inj(1, n), \qquad V_n = V, \qquad S_n = V^n$
- The algebra is given by
    - $lookup : X^V \longrightarrow X^{Loc}$
    - $update : X \longrightarrow X^{Loc \times V}$
    - $block : [L, X] \longrightarrow X^V$
    - subject to appropriate diagrams commuting
- $(Y^X)_n = [Inj, \mathrm{Set}](X - \times Inj(n, -), Y-)$
- $[X, Y]_n = [Inj, \mathrm{Set}](X-, Y(n + -))$

- See also work by Power (cotensoring models with comodels) and Staton (completeness via nominal sets)

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Refinement types

Refinement types

- Also known as predicate subtyping
- Assume we are given some simple types
  - $Nat, Loc, \ldots$
- But often we want to talk about refined versions of them
  - even natural numbers
  - odd natural numbers
  - open locations
  - closed locations
- Refinement types provide us with such a framework
  - "equipping your existing type system with suitable logic"

Refinement types

- Well-formedness of refinement types

$$\frac{}{\Gamma \vdash \sigma : Ref(\sigma)} \qquad \frac{\Gamma \vdash \phi : Ref(\sigma) \quad \Gamma, x : \phi \vdash P : wf}{\Gamma \vdash (x : \phi)P : Ref(\sigma)}$$

$$\frac{\Gamma \vdash \phi : Ref(\sigma_1) \quad \Gamma, x : \phi \vdash \psi : Ref(\sigma_2)}{\Gamma \vdash \Sigma_{x:\phi}\psi : Ref(\sigma_1 \times \sigma_2)} \qquad \frac{\Gamma \vdash \phi : Ref(\sigma) \quad \Gamma, x : \phi \vdash \psi : Ref(\tau)}{\Gamma \vdash \Pi_{x:\phi}\psi : Ref(\sigma \to \tau)}$$

- Examples of typing rules

$$\frac{\Gamma \vdash t : \phi \quad \Gamma \vdash P[t/x]}{\Gamma \vdash t : (x : \phi)P}$$

$$\frac{\Gamma, x : \phi \vdash t : \psi}{\Gamma \vdash \lambda x : \phi . t : \Pi_{x:\phi}\psi} \qquad \frac{\Gamma \vdash t_1 : \Pi_{x:\phi}\psi \quad \Gamma \vdash t_2 : \phi}{\Gamma \vdash t_1 t_2 : \psi[t_2/x]}$$

Refinement types

- Well-formedness of refinement types

$$\frac{}{\Gamma \vdash \sigma : \mathit{Ref}(\sigma)} \qquad \frac{\Gamma \vdash \phi : \mathit{Ref}(\sigma) \quad \Gamma, x : \phi \vdash P : \mathit{wf}}{\Gamma \vdash (x : \phi)P : \mathit{Ref}(\sigma)}$$

$$\frac{\Gamma \vdash \phi : \mathit{Ref}(\sigma_1) \quad \Gamma, x : \phi \vdash \psi : \mathit{Ref}(\sigma_2)}{\Gamma \vdash \Sigma_{x:\phi}\psi : \mathit{Ref}(\sigma_1 \times \sigma_2)} \qquad \frac{\Gamma \vdash \phi : \mathit{Ref}(\sigma) \quad \Gamma, x : \phi \vdash \psi : \mathit{Ref}(\tau)}{\Gamma \vdash \Pi_{x:\phi}\psi : \mathit{Ref}(\sigma \to \tau)}$$

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Refinement types

- Set-theoretic semantics (ala. Denney)
  - Interpret refinement type $\Gamma \vdash \phi : \mathit{Ref}(\sigma)$ as a family of PERs $[\![\Gamma]\!] \longrightarrow \mathit{PER}([\![\sigma]\!])$
  - other type constructors (sums,products) are interpreted straightforwardly
  - terms $\Gamma \vdash t : \phi$ are interpreted as $[\![\Gamma]\!] \longrightarrow \mathcal{P}([\![\sigma]\!])$ (subsets denoting the 'total realizers')

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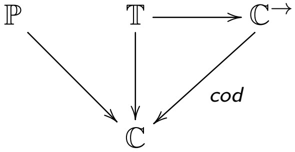

Refining global state

Refining global state

- We had the finite set of locations Loc
- Assume that we now have predicates $Open(Loc)$ and $Closed(Loc) = \neg Open(loc)$ on the locations Loc
- Conceptually they denote subsets of Loc
- We should only be able to read from and write to locations that are open
  - $lookup : X^V \longrightarrow X^{Open(Loc)}$
  - $update : X \longrightarrow X^{Open(Loc) \times V}$
- However, notice that this requires an a priori given collection of open locations

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Refining global state

- So we should also add operations for opening and closing locations
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- But we should be able to distinguish between computations able to use different locations
- We could take inspiration from the algebra for local state
  - work in the category Set^W
- However, we don't yet know what the appropriate non-discrete world category and the corresponding (monoidal) closed structure should be

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Refining global state (W-sorted theories)

- We don't know the definition in a single sorted theory
- So let's try to work in W-sorted algebraic theories
- A W-sorted algebraic theory consists of:
  - a set of sorts W (we think of them as worlds)
  - a small category $\mathcal{L}$ with countable products
  - an id. on objects countable-product preserving functor

    $$J : W^* \longrightarrow \mathcal{L}$$

  - *(where W^* has as objects words $w_0, \dots, w_n$ over W)*
- A model of a W-sorted theory is given by
  - a countable product preserving functor $M : \mathcal{L} \longrightarrow \mathrm{Set}$
- The forgetful functor $U : Mod(\mathcal{L}, \mathrm{Set}) \longrightarrow \mathrm{Set}^W$ again has a left adjoint F inducing a monad $T = UF$

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Refining global state (W-sorted theories)

- Let the worlds be $W = Bool^W$

- We have families of operations in the theory
  - $lookup_{w \in W, loc \in Open_w(Loc)} : w, ..., w \longrightarrow w$
  - $update_{w \in W, loc \in Open_w(Loc), v \in V} : w \longrightarrow w$
  - $open_{w \in W, loc \in Open_w(Loc)} : w \longrightarrow w[loc \mapsto \bot]$
  - $close_{w \in W, loc \in Closed_w(Loc)} : w \longrightarrow w[loc \mapsto \top]$
  - satisfying appropriate commuting diagrams

- Giving us the algebra
  - $lookup_{w \in W, loc \in Open_w(Loc)} : (X^V)_w \longrightarrow X_w$
  - $update_{w \in W, loc \in Open_w(Loc), v \in V} : X_w \longrightarrow X_w$
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Refining global state (W-sorted theories)

- So we have the algebra
  - $lookup_{w \in W, loc \in Open_w(Loc)} : (X^V)_w \longrightarrow X_w$
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- Inducing monad $TX_w = UFX_w = (\sum_{w' \in W} (S_{w'} \times X_{w'}))^{S_w}$
- With the unit $\eta_X : X \longrightarrow UFX$ of the adjunction given by:

  $\eta_{X,w}\, \gamma = \lambda s\,.\, \mathrm{inj}_w\,(s\,,\,\gamma)$
- And the counit $\varepsilon_A : FUA \longrightarrow A$ of the adjunction:

$$\varepsilon_{A,w} = (\coprod(S \times A_{w'}))^S \xrightarrow{(\coprod(S \times \overrightarrow{close}))^S} (\coprod(S \times A_{w^\top}))^S \xrightarrow{\cong} (S \times A_{w^\top})^S \xrightarrow{(\overrightarrow{write})^S} (A_{w^\top})^S \xrightarrow{\overrightarrow{read}} A_{w^\top} \xrightarrow{\overrightarrow{open}} A_w$$

- And the Kleisli extension is given by $(-)^* = U\varepsilon F$

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Another example of a straightforward theory

- Inspiration from McBride's work on file operations
- Take the simple set of worlds $W = Bool$
- We are interested in axiomatizing logging in to and logging off from some system
- Then we have the theory
  - $LogIn_{p \in Password} : true, false \longrightarrow false$
  - $DoSomething : true \longrightarrow true$
  - $LogOut : false \longrightarrow true$
- And the algebra
  - $LogIn_{p \in Password} : X_{true} \times X_{false} \longrightarrow X_{false}$
  - $DoSomething : X_{true} \longrightarrow X_{true}$
  - $LogOut : X_{false} \longrightarrow X_{true}$

What next?

- The W-sorted approach gave us the monad we were after
- Can we make it work naturally in the singlesorted case?
- Idea, try to give more general form to the operations in the algebra
  - $op_w : \prod_{o \in O_w} X_{\delta_o(w,o)} \longrightarrow \prod_{i \in I_w} X_{\delta_i(w,i)}$

  and in the theory
  - $op_w : \coprod_{o \in O_w} \{\delta_o(w,o)\} \longrightarrow \coprod_{i \in I_w} \{\delta_i(w,i)\}$
- But can't always define them uniformly in w, e.g.:

  $lookup_{[l_i \mapsto \bot]} : \coprod_{v \in V} \{[l_i \mapsto \bot]\} \longrightarrow 0$
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