

# String Diagrams *for* Free Monads

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Abstract nonsense  Practical programming



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*We want our reasoning principles to be:*

- **equational:** proofs are by chains of equalities between programs
- **expressive:** leveraging high-level properties of particular constructs

# free monads

A type of trees with nodes shaped by the functor  $f$  and leaves given by variables of the type  $a$ .

```
data Free f a
  = Var a
  | Op (f (Free f a))
```

Monadic bind given by substitution in leaves.

Sometimes denoted  $F^*$ .

# the problem

- Inductive data types (initial algebra) come with a low-level reasoning principle (structural induction).
- However, structural induction can be a bit awkward.

# the problem

- Inductive data types (initial algebra) come with a low-level reasoning principle (structural induction).
- However, structural induction can be a bit awkward.

Consider the **free monad transformer**:

```
newtype FreeT f m a = FreeT (m (Free (f . m) a))
```

Would **you** like to prove the monad laws using structural induction by hand?

Below, we prove that  $m \cdot \mu_K^K = M(\text{id} + \Sigma \mu_K^K) \cdot w$  and  $m \cdot K \mu^K = M(\text{id} + \Sigma K \mu^K) \cdot w$ , which means that both  $\mu_K^K$  and  $K \mu^K$  are coalgebra homomorphisms  $\mu_K^K, K \mu^K : \langle K^3, w \rangle \rightarrow \langle K^2, m \rangle$ . By uniqueness,  $\mu^K \cdot \mu_K^K = [m] \cdot \mu_K^K = [w] = [m] \cdot K \mu^K = \mu^K \cdot K \mu^K$ . Diagrammatically:

$$\begin{aligned}
 & m \cdot \mu_K^K \\
 = & \{ \text{definition of } m, \text{ universal property of } \mu^K \} \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2}) \cdot \alpha_K \cdot \alpha_K^{-1} \cdot M(\text{id}_K + g_{\Sigma K^2}) \cdot M(\text{id}_K + \text{id}_{\Sigma K^2} + \Sigma \mu_K^K) \\
 & \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{cancellation of } \alpha_K \text{ and } \alpha_K^{-1}, \text{ and naturality of } \alpha \} \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\text{id}_{M+\Sigma K} + g_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K} + \text{id}_{\Sigma K^2} + \Sigma \mu_K^K) \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2 + \Sigma K^2}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{naturality of flat} \} \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K} + g_{\Sigma K^2}) \\
 & \cdot M(\text{id}_{M+\Sigma K} + \text{id}_{\Sigma K^2} + \Sigma \mu_K^K) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2 + \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2 + \Sigma K^2}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{distributivity of composition over coproduct} \} \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K^2} + g_{\Sigma K^2}) \\
 & \cdot M(\text{id}_M + \Sigma \eta_K^K + \text{id}_{\Sigma K^2} + \Sigma \mu_K^K) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2 + \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2 + \Sigma K^2}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{naturality of } g_{\Sigma}, \text{ right unit for } K \} \\
 & M(\text{id} + \Sigma \mu_K^K) \cdot M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K^2} + g_{\Sigma K^2}) \\
 & \cdot M(\text{id}_M + \Sigma(\eta_K^K, \eta_K^K) + \Sigma \eta_K^K + \Sigma \text{id}_K) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2 + \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2 + \Sigma K^2}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{definition of } w \} \\
 & M(\text{id} + \Sigma \mu_K^K) \cdot w \\
 & \quad \quad \quad * * * \\
 & m \cdot K \mu^K \\
 = & \{ \text{definition of } m \} \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2}) \cdot \alpha_K \cdot K \mu^K \\
 = & \{ \text{naturality of } \alpha \} \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2}) \cdot M(\mu^K + \Sigma K \mu^K) \cdot \alpha_K
 \end{aligned}$$

$$\begin{aligned}
 & = \{ \text{properties of coproducts} \} \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_K + \Sigma K \mu^K) \cdot M(\mu^K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{universal property of } \mu^K \} \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_K + \Sigma K \mu^K) \cdot M(\alpha^{-1} + \text{id}_{\Sigma K^2}) \\
 & \cdot M(M(\text{id} + g_{\Sigma K^2}) + \text{id}_{\Sigma K^2}) \cdot M(M(\text{id} + \text{id}_{\Sigma K^2} + \Sigma \mu^K) + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\text{flat}_{M+\Sigma K^2, \Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot M(M(\alpha + \text{id}_{\Sigma K^2}) + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{cancellation of } \alpha \text{ and } \alpha^{-1} \} \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\text{id}_{M+\Sigma K} + \Sigma K \mu^K) \\
 & \cdot M(M(\text{id} + g_{\Sigma K^2}) + \text{id}_{\Sigma K^2}) \cdot M(M(\text{id} + \text{id}_{\Sigma K^2} + \Sigma \mu^K) + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\text{flat}_{M+\Sigma K^2, \Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot M(M(\alpha + \text{id}_{\Sigma K^2}) + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{naturality of flat} \} \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K} + \Sigma K \mu^K) \\
 & \cdot M(\text{id} + g_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot M(\text{id} + \text{id}_{\Sigma K} + \Sigma \mu^K + \text{id}_{\Sigma K^2}) \\
 & \cdot \text{flat}_{M+\Sigma K^2 + \Sigma K^2, \Sigma K^2} \cdot M(\text{flat}_{M+\Sigma K^2, \Sigma K^2} + \text{id}_{\Sigma K^2}) \\
 & \cdot M(M(\alpha + \text{id}_{\Sigma K^2}) + \text{id}_{\Sigma K^2}) \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{monad laws} \} \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K} \\
 & \cdot M(\text{id} + g_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot M(\text{id} + \text{id}_{\Sigma K} + \Sigma \mu^K + \text{id}_{\Sigma K^2}) \\
 & \cdot \text{flat}_{M+\Sigma K^2 + \Sigma K^2, \Sigma K^2} \cdot M(\text{flat}_{M+\Sigma K^2, \Sigma K^2} + \text{id}_{\Sigma K^2}) \\
 & \cdot M(M(\alpha + \text{id}_{\Sigma K^2}) + \text{id}_{\Sigma K^2}) \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{naturality of } \eta^K \} \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K} \\
 & \cdot M(\text{id} + g_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot M(\text{id} + \text{id}_{\Sigma K} + \Sigma \mu^K + \text{id}_{\Sigma K^2}) \\
 & \cdot \text{flat}_{M+\Sigma K^2 + \Sigma K^2, \Sigma K^2} \cdot M(\text{flat}_{M+\Sigma K^2, \Sigma K^2} + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{naturality of } g_{\Sigma} \} \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + g_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\text{id} + \Sigma \eta_K^K + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K} + \Sigma K \mu^K + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\text{id} + \text{id}_{\Sigma K} + \Sigma \mu^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2 + \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{right unit for } K, \text{ naturality of } g_{\Sigma} \} \\
 & M(\text{id} + \Sigma K \mu^K) \cdot M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + g_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\text{id} + \Sigma(\eta_K^K, \eta_K^K) + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2 + \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{definition of } w \} \\
 & M(\text{id} + \Sigma K \mu^K) \cdot w \\
 & \quad \quad \quad * * * \\
 & M(\text{id} + g_{\Sigma K^2}) \cdot M(\text{id} + g_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\text{id} + \Sigma \eta_K^K + \Sigma(\eta_K^K, \mu^K) + \Sigma K \mu^K) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2}
 \end{aligned}$$

# the problem

- Let's invent more abstract reasoning principles that respect the underlying structure!



# *string* diagrams

- A 2-dimensional notation for expressions
- Useful for presentation of more complicated expressions that involve a number of functors
- Some administrative equalities are **built-in**

# lookup

Consider the function

$$\text{lookup } k :: \forall a. \text{List (Key, a)} \rightarrow \text{Maybe } a$$

It is polymorphic in  $a$ . So, alternatively, we write

$$\text{lookup } k : \text{List} \circ (\text{Key}, -) \rightarrow \text{Maybe}$$

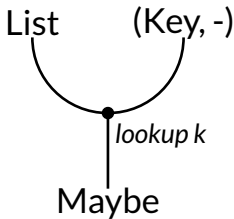
# lookup

Consider the function

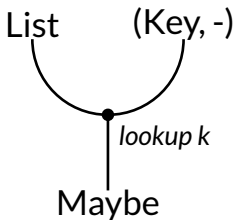
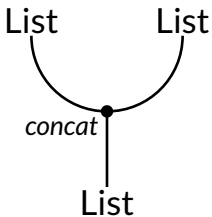
$\text{lookup } k :: \forall a. \text{List (Key, a)} \rightarrow \text{Maybe } a$

It is polymorphic in  $a$ . So, alternatively, we write

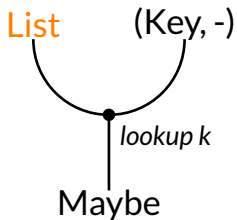
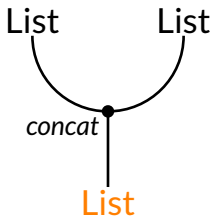
$\text{lookup } k : \text{List} \circ (\text{Key}, -) \rightarrow \text{Maybe}$



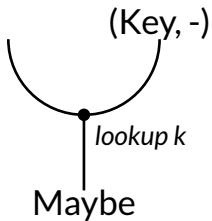
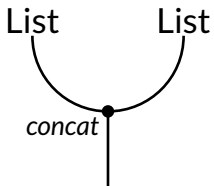
# composition



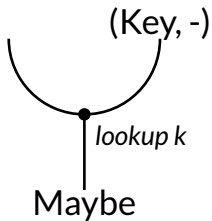
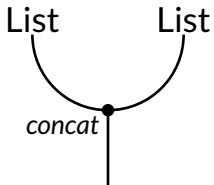
# composition



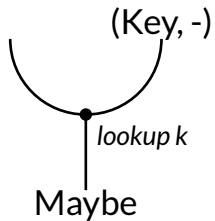
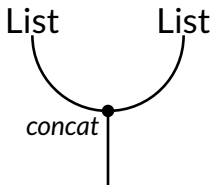
# composition



# composition

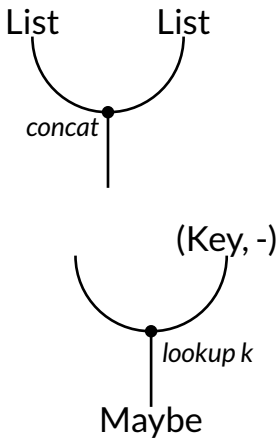


# composition

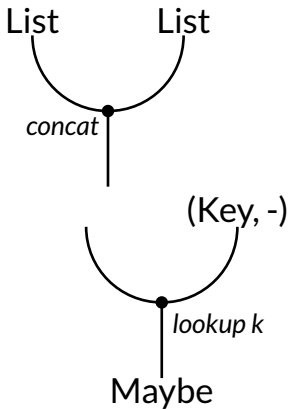




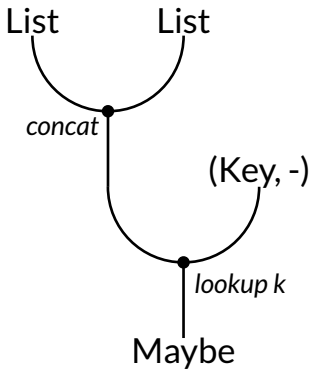
# composition



# composition

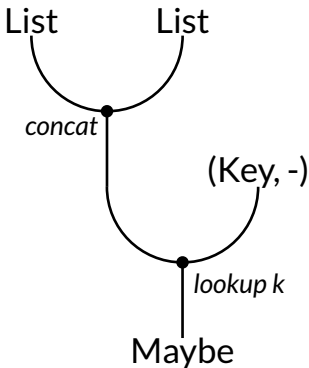


# composition



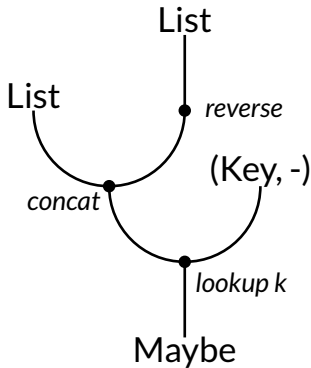
# composition

`lookup k . concat :: ∀a. List (List (Key, a)) → Maybe a`



# composition

`lookup k . concat . fmap reverse`  
`:: ∀a.List (List (Key, a)) → Maybe a`



# free monads, more abstractly

This definition is only one of many possible implementations:

```
data Free f a
  = Var a
  | Op (f (Free f a))
```

Shouldn't this be an interface?

```
class FreeMonad f a where ???
```

What are the **operations** and **laws** that characterise free monads?

# freeness, mathematically

$\text{emb} :: \forall a.f\ a \rightarrow \text{Free}\ f\ a$

$\text{emb} : F \rightarrow F^*$



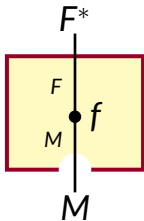
Generic operation that takes a single operation to a term.

# freeness, mathematically

`interp :: (Functor f, Monad m) =>  
 (forall x. f x -> m x) -> Free f a -> m a`

Given a monad  $M$  and  $f : F \rightarrow M$ ,  
there is a **monad morphism**

$$[f] : F^* \rightarrow M$$

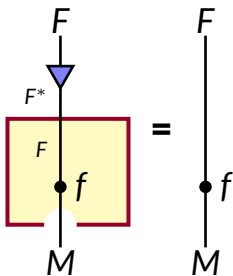




## cancellation

`interp f . emb = f`

$[f] \cdot \text{emb} = f$

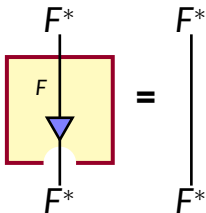


*Interpreting a term that consists of one operation is the same as interpreting that operation.*

# reflection

interp emb = id

[emb] = id

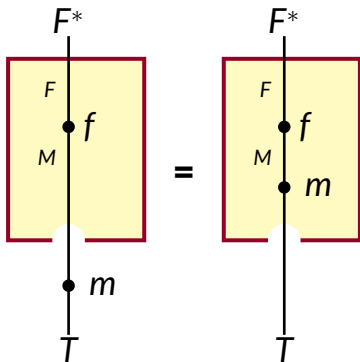


*Interpreting operations as syntax... preserves syntax.*

## fusion

$$m \cdot \text{interp } f = \text{interp } (m \cdot f)$$

$m \cdot [f] = [m \cdot f]$ , where  $m$  is a monad morphism



*Interpreting a term and then transforming semantics is the same as interpreting and transforming the term in one go.*

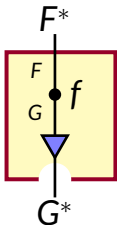
## example: renaming is functorial

Given functors  $F$  and  $G$ , and a natural transformation  $f : F \rightarrow G$ , we can define a natural transformation (*hoist*)

$$F^* \rightarrow G^*$$

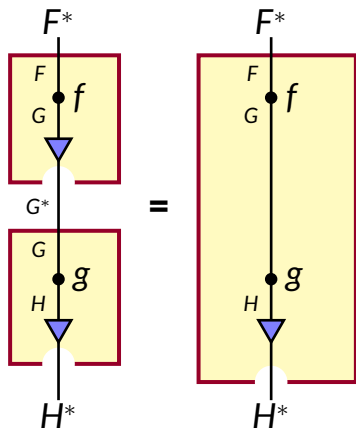
as follows:

$$[\text{emb} \cdot f]$$

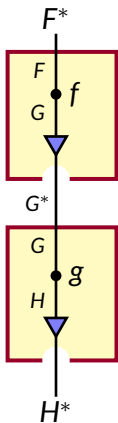


Given three functors  $F, G, H$ , and two natural transformations  $f : F \rightarrow G$  and  $g : G \rightarrow H$ , is the composition of hoists a hoist of composition? That is:

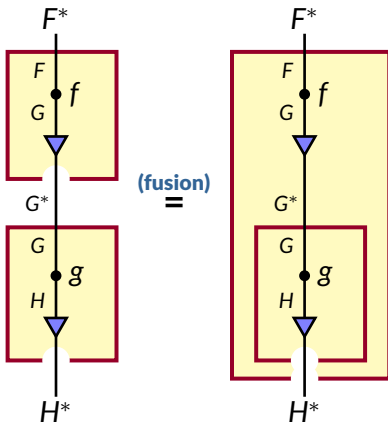
$$[\text{emb} \cdot g] \cdot [\text{emb} \cdot f] = [\text{emb} \cdot g \cdot f]$$



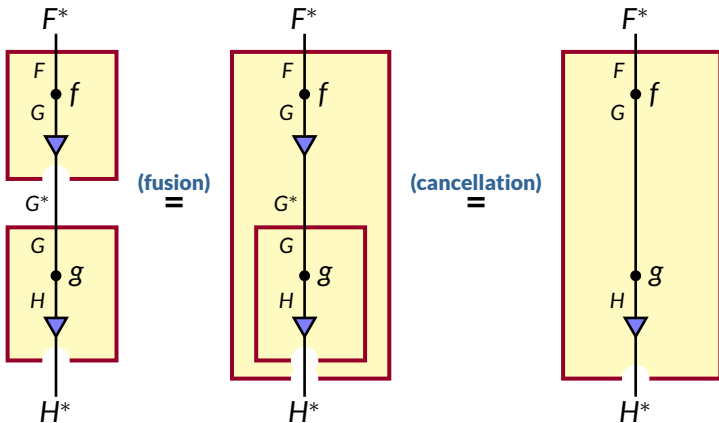
example: **renaming is functorial**



# example: renaming is functorial



# example: renaming is functorial





# going back to **FreeT**

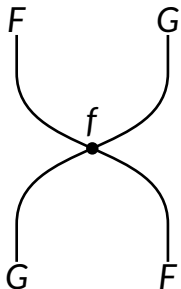
Sadly, what we have seen is not enough to show that **FreeT** is a monad.

Luckily, free monads in Haskell have other kinds of similar properties.

# distributable free monads

Given

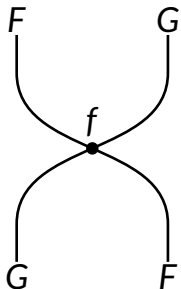
$$f : F \circ G \rightarrow G \circ F$$



# distributable free monads

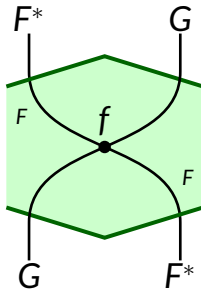
Given

$$f : F \circ G \rightarrow G \circ F$$



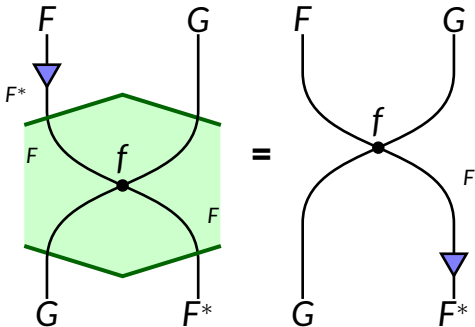
we get

$$\langle f \rangle : F^* \circ G \rightarrow G \circ F^*$$



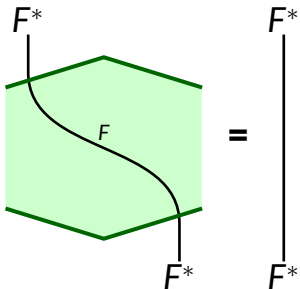
# cancellation

$$\langle f \rangle \cdot \text{emb} = G \text{emb} \cdot f$$



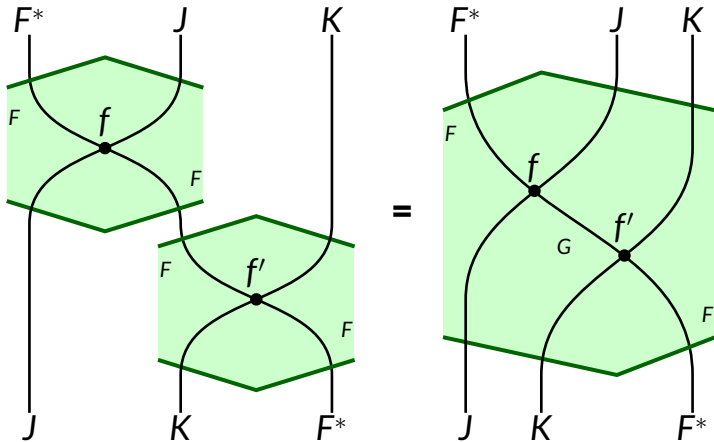
# reflection

$$\langle \text{id} \rangle = \text{id}$$



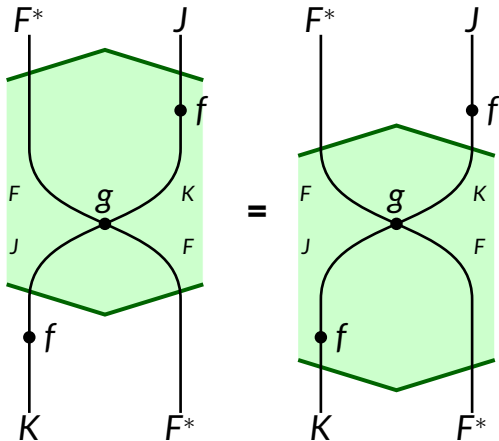
# fusion

$$J\langle f' \rangle \cdot \langle f \rangle K = \langle Jf' \cdot fK \rangle$$



# uniformity

$$fF^* \cdot \langle g \cdot F^*f \rangle = \langle fF \cdot g \rangle \cdot F^*f$$



# uniformity

Kind of

$$(fg)^*f = f(gf)^*$$

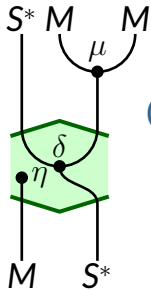
in Kleene algebra



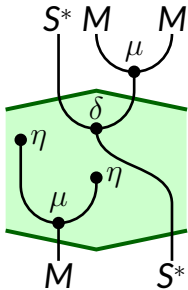
Below, we prove that  $m \cdot \mu_K^K = M(\text{id} + \Sigma \mu_K^K) \cdot w$  and  $m \cdot K\mu^K = M(\text{id} + \Sigma K\mu^K) \cdot w$ , which means that both  $\mu_K^K$  and  $K\mu^K$  are coalgebra homomorphisms  $\mu_K^K, K\mu^K : \langle K^3, w \rangle \rightarrow \langle K^2, m \rangle$ . By uniqueness,  $\mu^K \cdot \mu_K^K = [m] \cdot \mu_K^K = [w] = [m] \cdot K\mu^K = \mu^K \cdot K\mu^K$ . Diagrammatically:

$$\begin{aligned}
 & m \cdot \mu_K^K \\
 = & \{ \text{definition of } m, \text{ universal property of } \mu^K \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^3}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^3} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^3}) \cdot \alpha_K \cdot \alpha_K^{-1} \cdot M(\text{id}_K + g_{\Sigma K^2}) \cdot M(\text{id}_K + \text{id}_{\Sigma K^2} + \Sigma \mu_K^K) \\
 & \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{cancellation of } \alpha_K \text{ and } \alpha_K^{-1}, \text{ and naturality of } \alpha \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\text{id}_{M+\Sigma K} + g_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K} + \text{id}_{\Sigma K^2} + \Sigma \mu_K^K) \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2 + \Sigma K^3}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{naturality of flat} \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^3}) \cdot M(\text{id}_{M+\Sigma K} + g_{\Sigma K^2}) \\
 & \cdot M(\text{id}_{M+\Sigma K} + \text{id}_{\Sigma K^2} + \Sigma \mu_K^K) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2 + \Sigma K^3} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2 + \Sigma K^3}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{distributivity of composition over coproduct} \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id}_{M+\Sigma K^2} + g_{\Sigma K^2}) \\
 & \cdot M(\text{id}_M + \Sigma \eta_K^K + \text{id}_{\Sigma K^2} + \Sigma \mu_K^K) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2 + \Sigma K^3} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2 + \Sigma K^3}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{naturality of } g_{\Sigma}, \text{ right unit for } K \} \\
 & M(\text{id} + \Sigma \mu_K^K) \cdot M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id}_{M+\Sigma K^2} + g_{\Sigma K^2}) \\
 & \cdot M(\text{id}_M + \Sigma(\eta_K^K, \eta_K^K) + \Sigma \eta_K^K + \Sigma \text{id}_K) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2 + \Sigma K^3} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2 + \Sigma K^3}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{definition of } w \} \\
 & M(\text{id} + \Sigma \mu_K^K) \cdot w \\
 & \quad \quad \quad * * * \\
 & m \cdot K\mu^K \\
 = & \{ \text{definition of } m \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2}) \cdot \alpha_K \cdot K\mu^K \\
 = & \{ \text{naturality of } \alpha \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2}) \cdot M(\mu^K + \Sigma K\mu^K) \cdot \alpha_K
 \end{aligned}$$

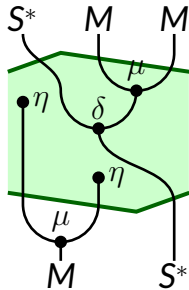
$$\begin{aligned}
 & = \{ \text{properties of coproducts} \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_K + \Sigma K\mu^K) \cdot M(\mu^K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{universal property of } \mu^K \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_K + \Sigma K\mu^K) \cdot M(\alpha^{-1} + \text{id}_{\Sigma K^2}) \\
 & \cdot M(M(\text{id} + g_{\Sigma K^2}) + \text{id}_{\Sigma K^2}) \cdot M(M(\text{id} + \text{id}_{\Sigma K^2} + \Sigma \mu^K) + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\text{flat}_{M+\Sigma K^2, \Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot M(M(\alpha + \text{id}_{\Sigma K^2}) + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{cancellation of } \alpha \text{ and } \alpha^{-1} \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{M+\Sigma K, \Sigma K^2} \\
 & \cdot M(\text{id}_{M+\Sigma K^2} + \Sigma K\mu^K) \\
 & \cdot M(M(\text{id} + g_{\Sigma K^2}) + \text{id}_{\Sigma K^2}) \cdot M(M(\text{id} + \text{id}_{\Sigma K^2} + \Sigma \mu^K) + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\text{flat}_{M+\Sigma K^2, \Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot M(M(\alpha + \text{id}_{\Sigma K^2}) + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{naturality of flat} \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K} + \Sigma K\mu^K) \\
 & \cdot M(\text{id} + g_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot M(\text{id} + \text{id}_{\Sigma K} + \Sigma \mu^K + \text{id}_{\Sigma K^2}) \\
 & \cdot \text{flat}_{M+\Sigma K^2 + \Sigma K^2, \Sigma K^2} \cdot M(\text{flat}_{M+\Sigma K^2, \Sigma K^2} + \text{id}_{\Sigma K^2}) \\
 & \cdot M(M(\alpha + \text{id}_{\Sigma K^2}) + \text{id}_{\Sigma K^2}) \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{monad laws} \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K} \\
 & \cdot M(\text{id} + g_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot M(\text{id} + \text{id}_{\Sigma K} + \Sigma \mu^K + \text{id}_{\Sigma K^2}) \\
 & \cdot \text{flat}_{M+\Sigma K^2 + \Sigma K^2, \Sigma K^2} \cdot M(\text{flat}_{M+\Sigma K^2, \Sigma K^2} + \text{id}_{\Sigma K^2}) \\
 & \cdot M(M(\alpha + \text{id}_{\Sigma K^2}) + \text{id}_{\Sigma K^2}) \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{naturality of } \eta^K \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K} \\
 & \cdot M(\text{id} + g_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot M(\text{id} + \text{id}_{\Sigma K} + \Sigma \mu^K + \text{id}_{\Sigma K^2}) \\
 & \cdot \text{flat}_{M+\Sigma K^2 + \Sigma K^2, \Sigma K^2} \cdot M(\text{flat}_{M+\Sigma K^2, \Sigma K^2} + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{naturality of } g_{\Sigma} \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K} \\
 & \cdot M(\text{id} + g_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot M(\text{id} + \text{id}_{\Sigma K} + \Sigma \mu^K + \text{id}_{\Sigma K^2}) \\
 & \cdot \text{flat}_{M+\Sigma K^2 + \Sigma K^2, \Sigma K^2} \cdot M(\alpha + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \\
 & \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{naturality of } g_{\Sigma} \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + \Sigma \eta_K^K + \text{id}_{\Sigma K^2}) \cdot M(\text{id}_{M+\Sigma K} \\
 & \cdot M(\text{id} + g_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot M(\text{id} + \text{id}_{\Sigma K} + \Sigma \mu^K + \text{id}_{\Sigma K^2}) \\
 & \cdot \text{flat}_{M+\Sigma K^2 + \Sigma K^2, \Sigma K^2} \cdot M(\alpha + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \\
 & \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2} \\
 = & \{ \text{properties of coproducts} \} \\
 & M(\text{id} + g_{\Sigma K^3}) \cdot M(\text{id} + g_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \\
 & \cdot M(\text{id} + \Sigma \eta_K^K + \Sigma(\eta_K^K, \mu^K) + \Sigma K\mu^K) \cdot \text{flat}_{M+\Sigma K} \\
 & \cdot M(\alpha + \text{id}_{\Sigma K^2} + \text{id}_{\Sigma K^2}) \cdot \text{flat}_{K+\Sigma K^2, \Sigma K^2} \cdot M(\alpha_K + \text{id}_{\Sigma K^2}) \cdot \alpha_{K^2}
 \end{aligned}$$



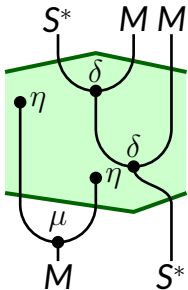
(unit)  
=



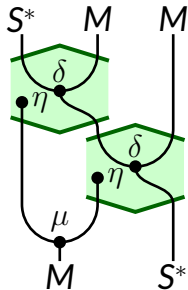
(uniformity)  
=



(associativity)  
=



(fusion)  
=



# more in the ICFP 2016 paper

- one more universal property (generalised fold)
- relationship between different universal properties
- a lot of further examples and an equational proof of a real-life theorem
- explanation where all these properties come from

